

Exactness of Semidefinite Relaxations for Nonlinear Optimization Problems with Underlying Graph Structure

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Polynomial Optimization

□ Polynomial Optimization:

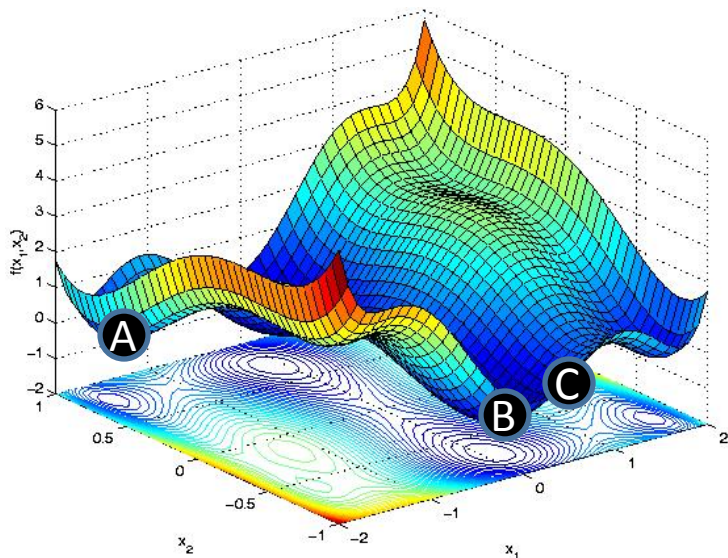
$$\begin{array}{ll}\min & x^T M x \\ \text{s.t.} & x_i^2 = 1, \quad i = 1, 2, \dots, n\end{array}$$



Special case: Combinatorial optimization and integer programming problems

Very hard to solve

□ Different types of solutions:



Point A: Local solution

Point B: Global solution

Point C: Near-global solution

Focus of our research

❖ **Approach:** Low-rank optimization, matrix completion, graph theory, convexification

Convexification

$$\begin{array}{ll} \min_{x \in \mathbb{C}^n} & x^H M_0 x \\ \text{s.t.} & x^H M_i x \leq a_i, \quad i = 1, 2, \dots, m \end{array}$$

$\text{trace}\{M_0 x x^H\}$



SDP relaxation

$$\begin{array}{ll} \min_{W \in \mathbb{H}^n} & \text{trace}\{M_0 W\} \\ \text{s.t.} & \text{trace}\{M_i W\} \leq a_i, \quad i = 1, 2, \dots, m \\ & W \succeq 0 \end{array}$$



Penalized SDP

$$\begin{array}{ll} \min_W & \text{trace}\{M_0 W\} + \lambda g(W) \\ \text{s.t.} & \text{trace}\{M_i W\} \leq a_i, \quad i = 1, 2, \dots, m \\ & W \succeq 0 \end{array}$$

❑ **Transformation:** Replace $x x^H$ with W .

❑ W is positive semidefinite and **rank 1**

❑ **Rank-1 SDP:** Recovery of a global solution x

❑ **Rank-1 penalized SDP:** Recovery of a near-global solution x

Research Problems

Arbitrary Real/Complex Polynomial Optimization



Conversion

$$\begin{aligned} \min_{x \in \mathbf{D}^n} \quad & x^H M_0 x \\ \text{s.t.} \quad & x^H M_i x \leq a_i, \quad i = 1, 2, \dots, m \end{aligned}$$



SDP/ Penalized SDP

$$\begin{aligned} \min_W \quad & \text{trace}\{M_0 W\} + \lambda g(W) \\ \text{s.t.} \quad & \text{trace}\{M_i W\} \leq a_i, \quad i = 1, 2, \dots, m \\ & W \succeq 0 \end{aligned}$$

How does structure make SDP relaxation exact?



Complexity analysis based on generalized weighted graph

Connection between sparsity and rank?



Proof of existence of low-rank solution using OS and treewidth

How to design penalized SDP?



Propose two methods to design penalty

Design scalable numerical algorithm?



Cheap iterations for large-scale problems

Power optimization problems



Finding near-global solutions using physics of power grids

Structured Optimization

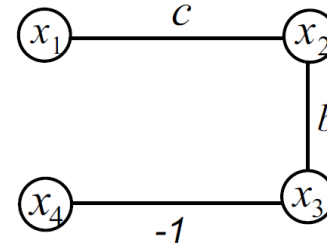
□ **Approach:** Map the structure into a *generalized weighted graph*.

$$\min_{x_1, x_2} x_1^4 + ax_2^2 + bx_1^2x_2 + cx_1x_2$$

$$\min_{x \in \mathbb{R}^4} x_3^2 + ax_2^2 + \underbrace{bx_2x_3} + \underbrace{cx_1x_2}$$

$$\text{s.t. } x_1^2 - \underbrace{x_3x_4} \leq 1$$

$$x_4^2 - 1 = 0$$

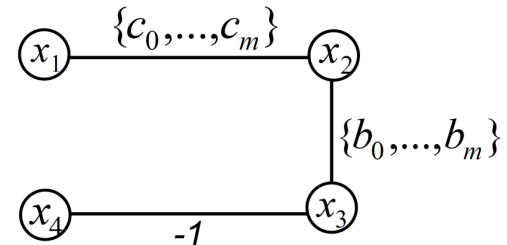


Due to structure, SDP is always exact.

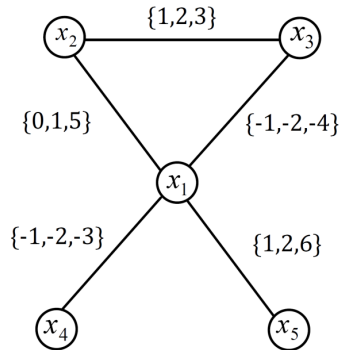
Generalized weighted graph:

$$\min_{x_1, x_2} x_1^4 + a_0x_2^2 + \underbrace{b_0x_1^2x_2} + c_0x_1x_2$$

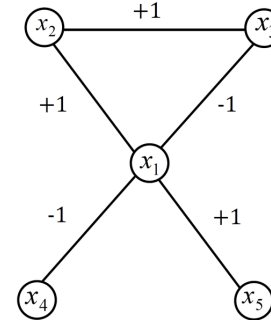
$$\text{s.t. } x_1^4 + a_ix_2^2 + \underbrace{b_ix_1^2x_2} + c_ix_1x_2 \leq \alpha_i, \quad i = 1, 2, \dots, m$$



Real-Valued Optimization



Sign assignment



Theorem

The SDP relaxation is exact if

$$\begin{aligned} \sigma_{ij} &\neq 0, & \forall (i, j) \in \mathcal{G} \\ \prod_{(i, j) \in \mathcal{O}_r} \sigma_{ij} &= (-1)^{|\mathcal{O}_r|}, & \forall r \in \{1, \dots, p\} \end{aligned}$$



Edge



Cycle

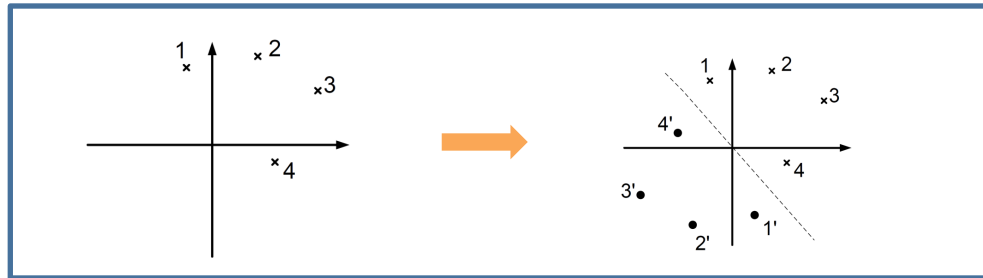
□ Special cases:

- ❖ **Positive optimization:** Bipartite graph
- ❖ **Negative optimization:** Arbitrary graph

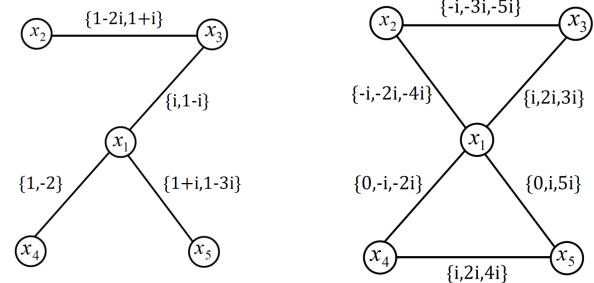
Interesting phenomena happen for complex optimization.

Complex-Valued Optimization

- ❑ **Real-valued case:** “ T ” is sign definite if T and $-T$ are separable in \mathbf{R} :
- ❑ **Complex-valued case:** “ T ” is sign definite if T and $-T$ are separable in \mathbf{R}^2 :



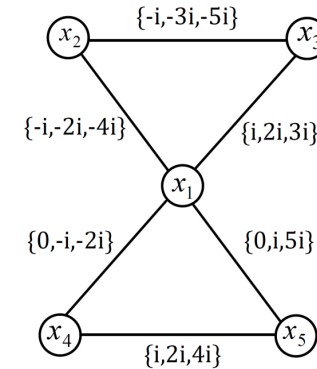
Theorem: SDP is exact for acyclic graphs with sign definite sets and certain cyclic graphs.



- ❑ The proposed conditions include several existing ones ([Kim and Kojima, 2003], [Padberg, 1989], [Bose, Gayme, Chandy, and Low, 2012], etc.).

Complex-Valued Optimization

- **Purely imaginary weights** (lossless power grid):



Theorem

Exact relaxation for weakly cyclic graphs with homogeneous weight sets.

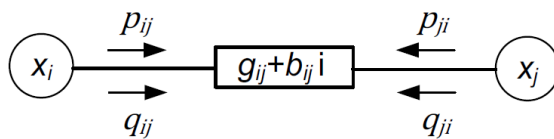
- Consider a real matrix \mathbf{M} :

$$\begin{aligned} \min_{\mathbf{x} \in \mathcal{C}^n} \quad & \mathbf{x}^* \mathbf{M} \mathbf{x} \\ \text{s.t.} \quad & |x_j| = 1, \quad j = 1, 2, \dots, m \end{aligned}$$

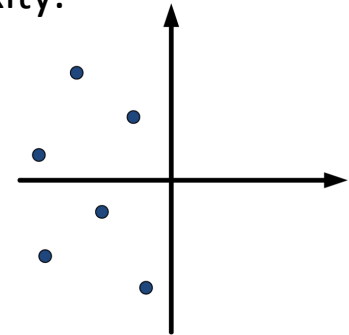
- Polynomial-time solvable for weakly-cyclic bipartite graphs.

Example

Example: Physics of power grids reduces computational complexity.



Coefficients of $x_i x_j$

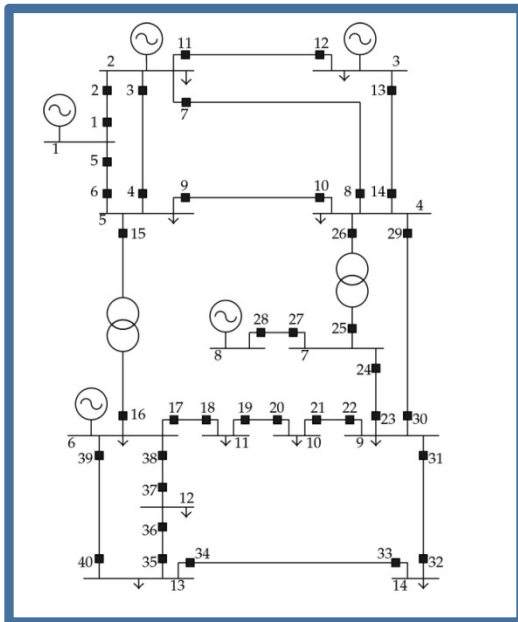
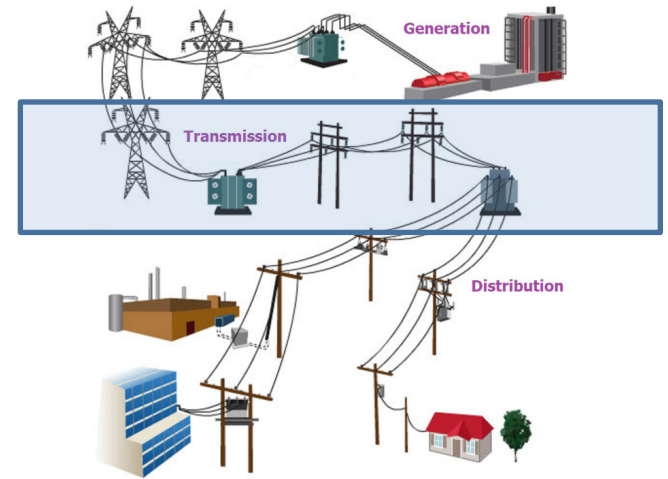


Sign definite due to passivity

Power Systems

❑ Power system:

- ❖ A large-scale system consisting of generators, loads, lines, etc.
- ❖ Used for generating, transporting and distributing electricity.



ISO, RTO, TSO



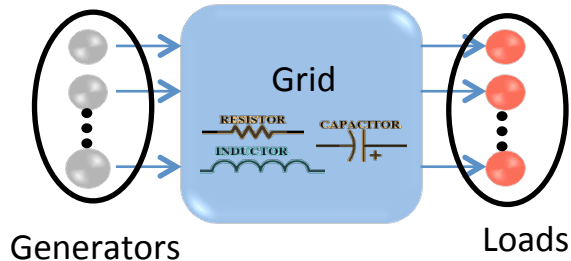
1. Optimal power flow (OPF)
2. Security-constrained OPF
3. State estimation
4. Network reconfiguration
5. Unit commitment
6. Dynamic energy management

NP-hard

(real-time operation and market)

Optimal Power Flow

Optimal Power Flow: Optimally match supply with demand

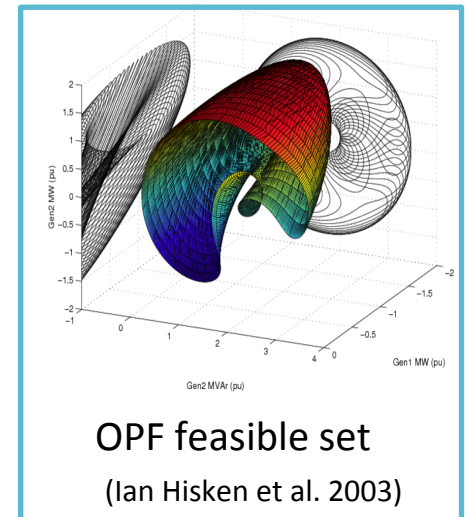


$$\begin{aligned} \min_{x \in \mathbb{C}^n} \quad & x^H M_0 x \\ \text{s.t.} \quad & x^H M_i x \leq a_i, \quad i = 1, 2, \dots, m \end{aligned}$$

Vector of complex voltages

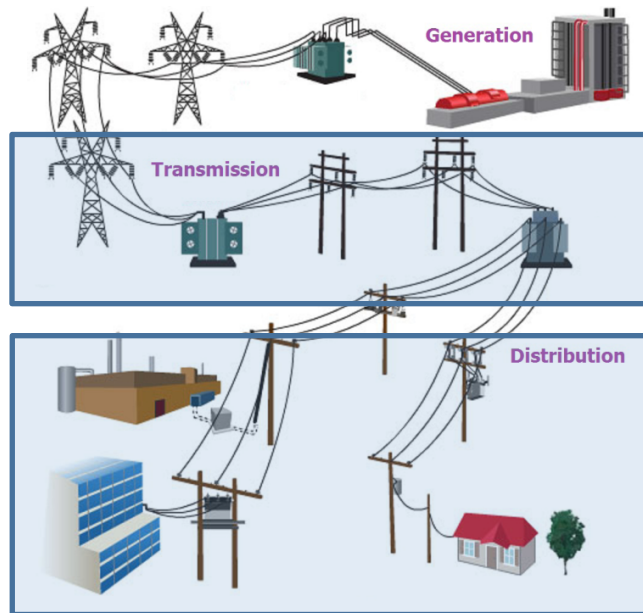
- ❑ **Real-time operation:** OPF is solved every 5-15 minutes.
- ❑ **Market:** Security-constrained unit-commitment OPF
- ❑ **Complexity:** Strongly NP-complete with long history since 1962.
- ❑ **Common practice:** Linearization
- ❑ **FERC and NETSS Study:** Annual cost of approximation > \$ 1 billion

A multi-billion critical system depends on optimization.



Exactness of Relaxation

- SDP is exact for IEEE benchmark examples and several real data sets.



cyclic



Theorem: Exact under positive LMPs with many transformers.

acyclic



Theorem: Exact under positive LMPs.

Physics of power networks (e.g., passivity) reduces computational complexity for power optimization problems.

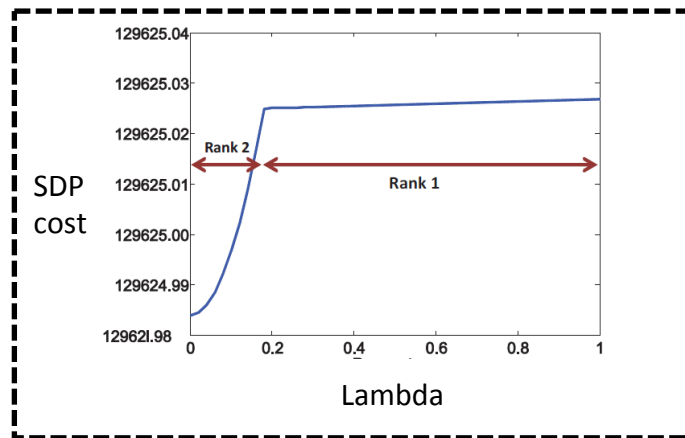
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2. S. Sojoudi and J. Lavaei, "Physics of Power Networks Makes Hard Optimization Problems Easy to Solve," PES 2012.
3. J. Lavaei and S. Low, "Zero Duality Gap in Optimal Power Flow Problem," IEEE Transactions on Power Systems, 2012.
4. J. Lavaei, D. Tse and B. Zhang, "Geometry of Power Flows and Optimization in Distribution Networks," IEEE Transactions on Power System, 2014.
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Example: Near-Global Solutions

Strategy: Penalize reactive loss over problematic lines

□ Modified IEEE 118-bus:

- ❖ 3 local solutions
- ❖ Costs: 129625, 177984, 195695



Case	TW	Cost	Guarantee	Time (sec)
Chow's 9 bus	2	5296.68	100%	≤ 5
IEEE 14 bus	2	8081.53	100%	≤ 5
IEEE 24 bus	4	63352.20	100%	≤ 5
IEEE 30 bus	3	576.89	100%	≤ 5
NE 39 bus	3	41864.40	99.994%	≤ 5
IEEE 57 bus	5	41737.78	100%	≤ 5
IEEE 118 bus	4	129660.81	99.995%	≤ 5
IEEE 300 bus	6	719725.10	99.998%	13.9
Polish 2383wp	23	1874322.65	99.316%	529
Polish 2736sp	23	1308270.20	99.970%	701
Polish 2737sop	23	777664.02	99.995%	675
Polish 2746wop	23	1208453.93	99.985%	801
Polish 2746wp	24	1632384.87	99.962 %	699
Polish 3012wp	24	2608918.45	99.188%	814
Polish 3120sp	24	2160800.42	99.073 %	910

Case	Minima	Cost	Guarantee
WB2	2	877.78	100%
WB3	2	417.25	100%
WB5	2	946.58	99.995%
WB5 Mod	3	1482.22	100%
LMBM3	5	5694.54	100%
LMBM3_50	2	5823.86	99.807%
case22loop	2	4538.80	100 %
case30loop	2	2863.06	100%
case30loop Mod	3	2861.88	100%
case39 Mod4	3	557.15	99.999%
case118 Mod1	3	129625.19	99.999%
case118 Mod2	2	85987.59	100 %
case300 Mod2	2	474643.46	99.996%

7000 simulations

Funding Acknowledgements

❑ **ONR YIP:** Graph-theoretic and low-rank optimization



❑ **DARPA YFA:** Near-Global Solutions of Non-convex Problems



❑ **NSF CAREER:** Control and optimization for power systems



❑ **NSF EPCN:** Contingency analysis for power systems



❑ **Google:** Numerical algorithms for nonlinear optimization



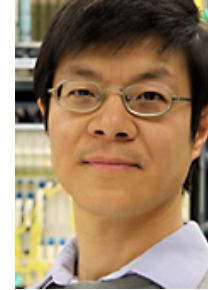
❑ **Siebel:** Computational methods for maximizing efficiency, reliability and resiliency of power systems



Collaborators

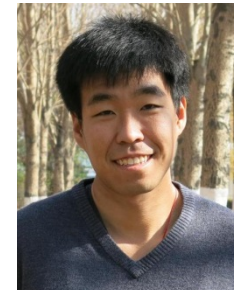
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- Steven Low
- Ross Baldick



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