

Scalable and Robust State Estimation from Abundant but Untrusted Data

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Abstract—Power system state estimation is an important problem in grid operation that has a long tradition of research since 1960s. Due to the nonconvexity of the problem, existing approaches based on local search methods are susceptible to spurious local minima, which could endanger the reliability of the system. In general, even in the absence of noise, it is challenging to provide a practical condition under which one can uniquely identify the global solution due to its NP-hardness. In this study, we propose a linear basis of representation that succinctly captures the topology of the network and enables an efficient two-stage estimation method in case the amount of measured data is not too low. Based on this framework, we propose an identifiability condition that numerically depicts the boundary where one can warrant efficient recovery of the unique global minimum. Furthermore, we develop a robustness metric called “mutual incoherence,” which underpins theoretical analysis of global recovery condition and statistical error bounds in the presence of both dense noise and bad data. The method demonstrates superior performance over existing methods in terms of both estimation accuracy and bad data robustness in an array of benchmark systems. Above all, it is scalable to large systems with more than 13,000 buses and can achieve accurate estimation within a minute.

Index Terms—Power system state estimation, statistical analysis, robust learning, smart grid

I. INTRODUCTION

Power system state estimation (SE) is conducted on a regular basis (e.g., every few minutes) to monitor the state of the grid by collecting and filtering a wealth of sensor data from transmission and distribution infrastructures [1], [2]. The state estimate presents system operators with essential information about the real-time operating status to improve situational awareness, make economic decisions, and take contingency actions in response to potential threats that could engender the grid reliability [3].

Due to the nonlinearity of the alternating-current (AC) grid physics, solving the set of power flow equations that arise from sensor measurements is known to be NP-hard for both transmission and distribution networks [4], [5]. As a result, there is a long tradition of studying this problem [2], [6]–[16]. At a high level, these methods are evaluated against multiple key criteria, including (i) accuracy (e.g., linearization/approximation of the nonlinear law of physics and its side effect on losing important information), (ii) robustness (e.g., to random/adversarial bad data, model mismatch, topological

errors), and (iii) scalability (i.e., computational/memory requirements to solve for large-scale systems). We provide a summary of the existing methods below, and refer the reader to [17] and [15] for a more comprehensive review.

A. Background and related work

The current practice in the power industry relies on a set of linearization and/or Newton’s methods that are originally developed in 1960s [2], [6], [7]. The Newton’s method has been employed to solve the nonlinear least square (NLS) SE and has quadratic convergence whenever the initial point is sufficiently close to the true solution [6]. However, the estimator is prone to outliers and sparse noise/errors, which can arise from sensor faults, topological errors [18]–[21], or adversarial attack [22]–[24]. To deal with large and sparse noise, one common approach is to perform bad data detection (BDD) on residual errors [25], [26]. This method relies on the statistical assumptions of the random errors (e.g., mean-zero and independent Gaussian distributions) and is only effective when the estimation from the Newton algorithm is close enough to the ground truth [2]. Alternatively, by redesigning the cost functions, robust estimators such as the least-absolute value (LAV) (a.k.a., ℓ_1 loss), the least median of squares, or Huber’s estimator have been employed [2], [8], [9], [27]–[30]. A major drawback of the above local search methods is the vulnerability to spurious local minima, which are those points that satisfy first- and second-order optimality conditions but are not the global minimum [30]–[32]. This is a major issue that can potentially cause a tremendous danger to the operation of the system, and it is difficult to distinguish a spurious local minimum that is consistent with data from the ground truth. Even though some recent works have shed light on the possibility of the non-existence of local minima in certain scenarios [33], the conditions are difficult to verify for SE [30].

Apart from local search algorithms mentioned above, several advanced optimization techniques have been proposed, such as particle swarm optimization [34], holomorphic embedding load flow method [35], homotopy continuation methods [36], feasible point pursuit [37], composite optimization [38], iterative mixed objective convex program [39], and algorithms for solving variational inequalities involving monotone operators [31]. A comprehensive review of these methods can be found in [15], [17].

The technique of convexification and semidefinite programming (SDP) relaxation is a powerful tool to tackle polynomial optimization, which arise from several areas such as graph

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theory, signal processing, and power systems [40]–[45]. Recently, the SDP relaxation technique has been applied to SE following its success for the optimal power flow problem [46], which demonstrate satisfactory numerical performance even in the presence of topological errors and bad data [11], [13], [16], [20], [47]. Also, [24] analyzes the vulnerabilities of AC SE against potential cyber attacks. While SDP relaxation is a promising approach with numerical success, this method requires that the solution to be rank-1 to recover the true state. Since the SDP relaxation is not always exact, one needs to add an extra rank penalty to the objective function (e.g., nuclear norm [11] or custom-designed norm [13], [16]), which would make the solution near-global optimal. The existing theoretical analyses of the SDP technique have not considered the bad data detection [13], [16]. Furthermore, the positive semidefinite constraint of the SDP technique limits the applicability of this method to large-scale problems, since the most common conic numerical algorithms scale on the order of $O(n^6)$, where n is the number of variables.

B. Contributions

We propose a method to solve large-scale AC SE with linear programming (LP) or quadratic programming (QP) that finds the correct state and is robust to sparse bad data, provided that the amount of measured data is relatively high. A new basis of representation is proposed, which is related to, but different from, the two dominant complex number representations used in power flow equations (i.e., polar coordinates and rectangular representation). This basis fully captures the properties of the power grid topology, which leads to efficient SE algorithms. Furthermore, we develop a SE identifiability condition to guarantee that no spurious local minimum exists in SE. We also perform theoretical analysis on the recovery condition of the true state in the presence of sparse bad data with statistical bounds on the estimation error.

The paper is organized as follows. The linear basis of representation is introduced in Sec. II-B, together with the measurement models and some key definitions to facilitate theoretical analysis. The two-stage estimator is introduced in Sec. III, whose performance is analyzed in Sec. IV. Sec. V includes numerical evaluations of the proposed methods on benchmark systems. Conclusion is drawn in Sec. VI. All proofs have been delegated to the appendix for the interested readers without interrupting the flow of the presentation.

II. POWER SYSTEM AC-MODEL

A. Notations

Vectors are shown by bold letters, and matrices are shown by bold and capital letters. Let x_i denote the i -th element of vector \mathbf{x} . We use \mathbb{R} and \mathbb{C} to show the sets of real and complex numbers. The set of indices $\{1, 2, \dots, m\}$ is denoted by $[m]$. The cardinality $|J|$ of a set J is the number of elements in the set. The support $\text{supp}(\mathbf{x})$ of a vector \mathbf{x} is the set of indices of the nonzero entries of \mathbf{x} . For a set $J \subseteq [m]$, we use $J^c = [m] \setminus J$ to denote its complement. We use \mathbf{A}_J to denote the submatrix formed by the rows of \mathbf{A} indexed by J . The symbol $(\cdot)^{\top}$ represents the transpose operator. We use $\langle \cdot \rangle$, $\text{Re}(\cdot)$ and $\text{Im}(\cdot)$ to denote the real part, imaginary part and trace of a scalar/matrix. The imaginary unit is denoted as i . The notations $\angle x$ and $|x|$ indicate the angle and magnitude of a complex scalar. For a convex function $g(\mathbf{x})$, we use $\partial g(\mathbf{x})$ to denote its subgradient. The inner product between two vectors is denoted by $\mathbf{h}^{\top} \mathbf{i}$. The notations $\|\mathbf{x}\|_1$, $\|\mathbf{x}\|_2$ and $\|\mathbf{x}\|_{\infty}$ represent the 1-norm, 2-norm and ∞ -norm of \mathbf{x} . We use \mathbb{E} to denote the expectation operator of a random variable.

$\text{Tr}(\cdot)$ to denote the real part, imaginary part and trace of a scalar/matrix. The imaginary unit is denoted as i . The notations $\angle x$ and $|x|$ indicate the angle and magnitude of a complex scalar. For a convex function $g(\mathbf{x})$, we use $\partial g(\mathbf{x})$ to denote its subgradient. The inner product between two vectors is denoted by $\mathbf{h}^{\top} \mathbf{i}$. The notations $\|\mathbf{x}\|_1$, $\|\mathbf{x}\|_2$ and $\|\mathbf{x}\|_{\infty}$ represent the 1-norm, 2-norm and ∞ -norm of \mathbf{x} . We use \mathbb{E} to denote the expectation operator of a random variable.

B. Power system modeling

We model the electric grid as a graph $G := (\mathcal{N}, \mathcal{L}, g)$, where $\mathcal{N} := [n_b]$ and $\mathcal{L} := [n_l]$ represent its set of buses and branches. Each branch $\ell \in \mathcal{L}$ that connects bus k and bus j is characterized by the branch admittance $y_{\ell} = g_{\ell} + ib_{\ell}$ and the shunt admittance $y_{\ell}^{\text{sh}} = g_{\ell}^{\text{sh}} + ib_{\ell}^{\text{sh}}$, where g_{ℓ} (resp., g_{ℓ}^{sh}) and b_{ℓ} (resp., b_{ℓ}^{sh}) denote the (shunt) conductance and susceptance, respectively. Since $g_{\ell}^{\text{sh}} = b_{\ell}^{\text{sh}} = 0$ in practice, we set it to zero in the subsequent description. In addition, to avoid duplicate definitions, each line $\ell := (k, j)$ is assigned with a unique direction from bus k (i.e., from end, given by $f(\ell) := k$) to bus j (i.e., to end, given by $t(\ell) := j$). We also use $\ell : \{k, j\}$ to denote a line ℓ with the direction of either (k, j) or (j, k) .

The power system state is described by the complex voltage vector $\mathbf{V} = [v_1, \dots, v_{n_b}]^{\top} \in \mathbb{C}^{n_b}$, where $v_k \in \mathbb{C}$ is the complex voltage at bus $k \in \mathcal{N}$ with magnitude $|v_k|$ and phase $\theta_k := \angle v_k$. Given the complex voltages, by Ohm's law, the complex current injected into line $\ell : \{k, j\}$ at bus k is given by:

$$i_{kj} = y_{\ell}(v_k - v_j) + \frac{i}{2} b_{\ell}^{\text{sh}} v_k.$$

Defining $\theta_{kj} := \theta_k - \theta_j$, one can write the power flow from bus k to bus j as

$$\begin{aligned} p_{kj}^{(\cdot)} &= |v_k|^2 g_{\ell} - |v_k||v_j|(g_{\ell} \cos \theta_{kj} - b_{\ell} \sin \theta_{kj}), \\ q_{kj}^{(\cdot)} &= |v_k|^2 (b_{\ell} + \frac{1}{2} b_{\ell}^{\text{sh}}) + |v_k||v_j|(b_{\ell} \cos \theta_{kj} - g_{\ell} \sin \theta_{kj}), \end{aligned}$$

and active (reactive) power injections at bus k ,

$$p_k = \sum_{\ell: f(\ell)=k} p_{kj}^{(\cdot)}, \quad q_k = \sum_{\ell: f(\ell)=k} q_{kj}^{(\cdot)}. \quad (1)$$

The above formulas are based on polar coordinates of complex voltages, where measurements are nonlinear functions of voltage magnitudes and phases. Another popular representation uses rectangular coordinates of complex numbers, where measurements are expressed as quadratic functions of the real and imaginary parts of voltages (see [48, Chap. 1] for more details).

C. Linear basis of representation

In this paper, we introduce a new basis of representation, where measurements can be expressed as *linear combinations* of the quantities derived from bus voltages. Specifically, for a given system G , we introduce two groups of variables:

- 1) voltage magnitude square, $x_k^{\text{mg}} := |v_k|^2$, for each bus $k \in \mathcal{N}$, and
- 2) real and imaginary parts of complex products, denoted as $x_{\ell}^{\text{re}} := \text{Re}(v_i v_j)$ and $x_{\ell}^{\text{im}} := \text{Im}(v_i v_j)$, respectively, for

each line $\ell = (i, j)$. Note that there is only one set of variables x^{re} and x^{im} for each line.

Using this representation, we can re-derive various types of power and voltage measurements (without noise) as follows:

Voltage magnitude square: The voltage magnitude square at bus $k \geq N$ is simply x_k^{mg} by definition.

Branch power flows: For each line $\ell = (i, j)$, the real and reactive power flows from bus i to bus j and in the reverse direction are given by:

$$\begin{aligned} p_{ij}^{(\cdot)} &= g \cdot x_i^{\text{mg}} \cdot g \cdot x^{\text{re}} + b \cdot x^{\text{im}} \\ q_{ij}^{(\cdot)} &= (b \cdot + \frac{1}{2} b^{\text{sh}}) x_i^{\text{mg}} + b \cdot x^{\text{re}} + g \cdot x^{\text{im}} \\ p_{ji}^{(\cdot)} &= g \cdot x_j^{\text{mg}} \cdot g \cdot x^{\text{re}} + b \cdot x^{\text{im}} \\ q_{ji}^{(\cdot)} &= (b \cdot + \frac{1}{2} b^{\text{sh}}) x_j^{\text{mg}} + b \cdot x^{\text{re}} + g \cdot x^{\text{im}} \end{aligned}$$

Nodal power injection: The power injection at bus node k consists of real and reactive powers, where:

$$\begin{aligned} p_k &= \sum_{k2} g \cdot x_k^{\text{mg}} \cdot g \cdot x^{\text{re}} + b \cdot x^{\text{im}} \\ q_k &= \sum_{k2} (b \cdot + \frac{1}{2} b^{\text{sh}}) x_k^{\text{mg}} + b \cdot x^{\text{re}} + g \cdot x^{\text{im}} \end{aligned}$$

where \sum_{k2} is the sum over all lines $\ell \geq L$ that are connected to k , $\sum_{f(\cdot)=k}$ is the sum over all lines ℓ where $f(\ell) = k$, and similarly, $\sum_{t(\cdot)=k}$ is the sum over all lines ℓ where $t(\ell) = k$. Equivalently, we can use (1) to combine the branch power flows defined above.

Thus, each customary measurement in power systems that belongs to one of the above *measurement types* can be represented by a linear function¹:

$$m_i(\mathbf{x}) = \mathbf{a}_i^T \mathbf{x}_1, \quad (2)$$

where $\mathbf{a}_i \geq R^{n_x}$ is the vector for the i -th noiseless measurement and $\mathbf{x}_1 = (f x_k^{\text{mg}}, g x^{\text{re}}, x^{\text{im}}) \geq R^{n_x}$ is the regression vector. By collecting all the sensor measurements in a vector $\mathbf{m} \geq R^{n_m}$, we have

$$\mathbf{m} = \mathbf{A} \mathbf{x}_1, \quad (3)$$

where $\mathbf{A} \geq R^{n_m \times n_x}$ is the sensing matrix with rows \mathbf{a}_i^T for $i \geq [n_m]$. Fig. 1 illustrates the sensing equation (3) for a simple 3-bus system.

It is worth mentioning that the linear basis introduced above is different from DC modeling of measurements, because the expression is *exact* for the AC model. This parametrization is inspired by the semidefinite relaxation approach for power system optimization [11], [13], [16], [20], [47], and it efficiently exploits the sparsity of the network (more on this in Sec. IV-A).

¹It is straightforward to include linear PMU measurements in our analysis as well using the relation $\tan \theta_{ij} = x_{ij}^{\text{im}} / x_{ij}^{\text{re}}$ for each line $\ell = (i, j)$, assuming we have a pair of PMUs on each end of a branch.

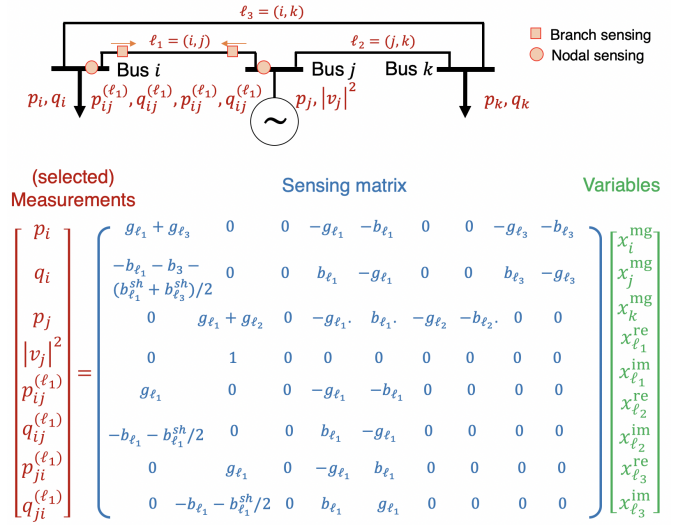


Fig. 1: Illustration of the sensing equation (3) for a 3-bus system. A selected set of measurements are considered, namely nodal injections at buses i and j , voltage magnitude square at bus j , and branch power flows along line $\ell_1 = (i, j)$. Note that one can choose the set of regression variables based on the availability of measurements, as long as each measurement can be fully represented by the chosen set of variables. For instance, we can omit the variables $x_{\ell_3}^{\text{re}} := \langle v_i v_k \rangle$ and $x_{\ell_3}^{\text{im}} := \langle v_i v_k \rangle$ by simultaneously excluding measurements $p_{ij}^{(\ell_3)}, q_{ij}^{(\ell_3)}, p_{ji}^{(\ell_3)}, q_{ji}^{(\ell_3)}$ and p_i, q_i, p_k, q_k , since they all rely on the omitted variables.

D. Measurement model

To perform SE, the supervisory control and data acquisition (SCADA) system collects measurements on power flows and complex voltages at key locations instrumented with sensors. This process is subject to both ubiquitous sensor noise and randomly occurring sensor faults. We consider the measurement model as follows:

$$\mathbf{y} = \mathbf{A} \mathbf{x}_1 + \mathbf{v} + \mathbf{b}_1, \quad (4)$$

where $\mathbf{A} \geq R^{n_m \times n_x}$ and $\mathbf{x}_1 \geq R^{n_x}$ are the sensing matrix and the true regression vector in (3), $\mathbf{v} \geq R^{n_m}$ denotes random noise, and $\mathbf{b}_1 \geq R^{n_m}$ is the bad data error that accounts for sensor failures or adversarial attacks [24]. Let $\mathcal{J} := \text{supp}(\mathbf{b}) \subseteq [n_m]$ denote the support of the bad data \mathbf{b} . We introduce the following properties to characterize the sensing matrix \mathbf{A} .

Definition 1 (Lower eigenvalue). Let $\mathbf{Q}_{\mathcal{J}} := \mathbf{A}_{\mathcal{J}}^T \mathbf{I}_{\mathcal{J}}$, where $\mathbf{I}_{\mathcal{J}}$ consists of the \mathcal{J} rows of the identity matrix $\mathbf{I} \geq R^{n_m \times n_m}$, and let $\mathbf{A}_{\mathcal{J}^c}$ be the submatrix of \mathbf{A} with rows indexed by \mathcal{J}^c . Then, the lower eigenvalue $C_{\min}(\mathcal{J})$ for a given corruption support \mathcal{J} is defined as the lower bound:

$$\min \lambda_{\min} \mathbf{Q}_{\mathcal{J}}^T \mathbf{Q}_{\mathcal{J}}, \lambda_{\min} \mathbf{A}_{\mathcal{J}^c}^T \mathbf{A}_{\mathcal{J}^c}, \quad (5)$$

where $\lambda_{\min}(\mathbf{X})$ denotes the smallest eigenvalue of \mathbf{X} .

The value $C_{\min}(\mathcal{J})$ characterizes the influence of bad data

on the identifiability of \mathbf{x}_1 . If $C_{\min}(\mathcal{J})$ is strictly positive, and one can accurately detect the support of bad data (a.k.a., support recovery), then it would be possible to obtain a good estimation of \mathbf{x}_1 with only the clean data in \mathcal{J}^c . Typically, the bad data due to sensor faults are randomly located, so if only a small amount of sensors are grossly corrupted (i.e., $j\mathcal{J} < j\mathcal{J}^c$), then the first term in (5) will be smaller than the second term. As we will see in Sec. IV, the first term is relevant for the case with dense noise \mathbf{w}_1 .

The next property turns out to be critical for BDD.

Definition 2 (Mutual incoherence). *Given a set $\mathcal{J} \subseteq [m]$ and its complement $\mathcal{J}^c := [m] \setminus \mathcal{J}$, let the pseudoinverse of $\mathbf{A}_{\mathcal{J}^c}$ be denoted as $\mathbf{A}_{\mathcal{J}^c}^+ = (\mathbf{A}_{\mathcal{J}^c}^T \mathbf{A}_{\mathcal{J}^c})^{-1} \mathbf{A}_{\mathcal{J}^c}^T$. Then, the mutual incoherence parameter $\rho(\mathcal{J})$ is defined to be:*

$$\rho(\mathcal{J}) = k \mathbf{A}_{\mathcal{J}^c}^+ \mathbf{A}_{\mathcal{J}}^T k_1,$$

where k k_1 denotes the matrix infinity norm (i.e., the maximum absolute column sum of the matrix).

The name ‘‘mutual incoherence’’ originates from the compressed sensing literature [49]–[52]. However, our definition is different because it measures the alignment of the sensing directions of the corrupted measurements (i.e., $\mathbf{A}_{\mathcal{J}}$) with those of the clean data (i.e., $\mathbf{A}_{\mathcal{J}^c}$). If these directions are misaligned (a.k.a., incoherent), then the value $\rho(\mathcal{J})$ is low and therefore it is likely to uncover the support of bad data. In general, the smaller the number of bad data measurement is, the more likely that $\rho(\mathcal{J})$ is small. As our analysis will show, if this value is strictly less than 1, then we can provably recover the support of the bad data.

Because the sensor data are of different types and scales, we make a normalization assumption.

Definition 3 (Measurement normalization). *Each row of \mathbf{A} is normalized as*

$$k \mathbf{a}_i k_2^2 = 1, \quad \forall i \in [n_m] \quad (6)$$

where \mathbf{a}_i is the i -th row of \mathbf{A} .

This condition is straightforward to implement in practice, since one can arbitrarily rescale the given coefficients of each measurement equation. This is also known as preconditioning, which assists with both the numerical stability and the statistical performance of regression.

III. TWO-STAGE STATE ESTIMATION

This section describes the proposed two-stage state estimation method, where both stages are linear or quadratic regression problems.

A. Stage 1: Estimation of \mathbf{x}_1

In the first stage, the goal is to estimate \mathbf{x}_1 from a set of noisy and corrupted measurements \mathbf{y} . We consider two cases separately. In the first case, the dense noise is negligible, i.e., $\mathbf{w}_1 = \mathbf{0}$ and we only need to consider the sparse measurement corruption \mathbf{b} . Under some conditions to be specified in Sec. IV, one can exactly recover the underlying vector \mathbf{x}_1 .

Case 1: Sparse corruption but no dense noise (i.e., $\mathbf{w} = \mathbf{0}$)

In this case, the measurements are given by $\mathbf{y} = \mathbf{A}\mathbf{x}_1 + \mathbf{b}_1$. To estimate \mathbf{x}_1 , we solve the following program:

$$\min_{\mathbf{x} \in \mathbb{R}^{n_x}; \mathbf{b} \in \mathbb{R}^{n_m}} k \mathbf{b} k_1, \quad \text{subject to } \mathbf{A}\mathbf{x} + \mathbf{b} = \mathbf{y}. \quad (\text{S1-L1})$$

Briefly, if the lower eigenvalue is bounded away from 0 (i.e., $C_{\min}(\mathcal{J}) > 0$) and the mutual incoherence is less than 1 (i.e., $\rho(\mathcal{J}) < 1$), then we can faithfully recover \mathbf{x}_1 and \mathbf{b}_1 from the above program.

Case 2: Sparse corruption and dense noise

In this case, the dense noise cannot be ignored, and the measurements are given by (4). We perform the estimation by solving the following LASSO-style optimization:

$$\min_{\mathbf{x} \in \mathbb{R}^{n_x}; \mathbf{b} \in \mathbb{R}^{n_m}} \frac{1}{2n_m} k \mathbf{y} - \mathbf{A}\mathbf{x} - \mathbf{b} k_2^2 + \lambda k \mathbf{b} k_1, \quad (\text{S1-LASSO})$$

where $\lambda > 0$ is the regularization coefficient. Due to the existence of dense noise, it is no longer possible to exactly recover the true \mathbf{x}_1 ; however, if the magnitudes of the dense noise are small, then we can still have good statistical bounds on the estimation error.

To detect and remove bad data, we first estimate the linear basis by solving either (S1-L1) or (S1-LASSO). This automatically produces a bad data vector estimation \mathbf{b} . If the value of any entry of \mathbf{b} is larger than a threshold, namely 0.1 in the experiments, then we classify the corresponding measurement as ‘‘bad data.’’ If there is a topology error, then the bad data tend to be localized on a line or a bus. We can observe the topological distribution of bad data vector to determine if there are such occurrences.

B. Stage 2: Recovery of \mathbf{v}

The goal of the second stage is to recover the underlying system voltage \mathbf{v} from the estimation $\hat{\mathbf{x}}$ from stage 1. First, we transform $\hat{\mathbf{x}}$ into estimations of voltage magnitudes and phase differences:

The voltage magnitude at each bus $k \in N$ is estimated as $\hat{v}_k = \hat{x}_k^{\text{mg}}$;

The phase difference along each line $\ell = (i, j)$ is estimated as $\hat{\theta}_{ij} = \arctan(\hat{x}_{ij}^{\text{im}} / \hat{x}_{ij}^{\text{re}})$.

To obtain the phase estimation at each bus, we solve the least-squares problem

$$\hat{\boldsymbol{\theta}} = \arg \min_{\boldsymbol{\theta} \in \mathbb{R}^{n_b}} \sum_{\ell=(i,j)} (\theta_i - \theta_j - \hat{\theta}_{ij})^2, \quad (\text{S2-}\theta)$$

which has a closed-form solution. To delve into this, let Δ be a collection of $\hat{\theta}_{ij}$, and $\mathbf{L} \in \mathbb{R}^{n_b \times n_b}$ be a sparse matrix with $L(\ell, i) := 1$ and $L(\ell, j) := -1$ for each line $\ell = (i, j)$ and zero elsewhere. Then, the solution for (S2- θ) is given by:

$$\hat{\boldsymbol{\theta}} = (\mathbf{L}^T \mathbf{L})^{-1} \mathbf{L}^T \Delta. \quad (7)$$

Finally, we can reconstruct $\hat{\mathbf{v}}$ by definition:

$$\hat{v}_k = \hat{v}_k e^{i \hat{\theta}_k}, \quad k \in N. \quad (8)$$

If the regression vector from stage 1 is exact, i.e., $\hat{x} = x$, then we can accurately recover the system state. Even if the \hat{x} is not exact, the second stage estimator (S2) has strong properties to control the estimation error and ensure that the errors in \hat{x}_{ij} will not propagate along the branches.

IV. THEORETICAL ANALYSIS

This section presents several theoretical analysis for the proposed framework. First, a condition for AC SE identifiability is presented. Then, we discuss the conditions under which accurate recovery of the true state is guaranteed. Furthermore, we present a novel statistical analysis of the recovery condition using concentration bounds.

A. Identifiability condition

Due to the nonconvexity of NLS, the existence of spurious local minima in SE is well-recognized, which makes it difficult to analyze whether the true state can be uniquely identified based on a given set of clean measurements (i.e., $b = 0$). Because SE can be formulated as a quadratic sensing problem, the results from the low-rank compressed sensing community seem to be directly applicable, which rely on a condition called restricted isometry property (RIP) (e.g., see [33], [53], [54]). The main result from this line of research indicates that if RIP of the sensing system is small enough, then every local minimum is also a global minimum [33], [54]. However, numerical results indicate that the condition is often too stringent to be satisfied for SE. It is also possible to characterize an “essentially strongly convex region” around the true solution, where any initial point converges to the true solution by local search [32], or to delineate a recovery region where the rank penalty leads to exact rank-1 solution [13], [16]. However, they all depend on the location of the true solution and the condition is hard to check numerically. The following theorem provides a condition similar to the DC-approximation results but for the AC SE. Without loss of generality, assume that the power network is connected.

Theorem 4. In the absence of noise (i.e., $b = 0$), one can uniquely identify the true state of the power grid if there exists a spanning tree T_{span} such that the set of measurements on the spanning tree (not including nodal injections) forms a matrix A_{span} that has full column rank (i.e., the null space of A_{span} is zero).

We do not include nodal injections in order to avoid dealing with variables on lines that do not have any measurements. For a given set of clean measurements, if the set satisfies the above identifiability condition, then there is one and only one global optimum for the Stage 1 estimator. Furthermore, this global optimum can be used in Stage 2 to recover the true state of the system.

Theorem 4 is inspired by the literature on false data injection attacks (FDIA) [22], [24], but the main difference is that we work through a completely different basis of variables. In FDIA, usually one works with the complex voltage directly. In that case, it is cumbersome to derive a necessary/sufficient condition for identifiability, because it is often algorithmically

(a) full network

(b) spanning tree

Fig. 2: Network topology of the IEEE 14 bus system. A sufficient condition for identifiability is the existence of a spanning tree T_{span} (shown on the right) where the measurements form a nonsingular sensing matrix A_{span} .

TABLE I: Empirical evaluation of the identifiability condition (Y/N) in Theorem 4. Measurement sets include: M1: $f_j v_k j g_{k2N}$ and $f p_{ij}^{(c)}; q_{ij}^{(c)} g_{2T_{\text{span}}}$; M2: $f_j v_k j g_{k2N}$ and $f p_{ij}^{(c)}; p_{ji}^{(c)} g_{2T_{\text{span}}}$; M3: $f_j v_k j g_{k2N}$ and $f q_{ij}^{(c)}; q_{ji}^{(c)} g_{2T_{\text{span}}}$; M4: $f p_{ij}^{(c)}; p_{ji}^{(c)}; q_{ij}^{(c)} g_{2T_{\text{span}}}$; M5: $f p_{ij}^{(c)}; q_{ij}^{(c)}; q_{ji}^{(c)} g_{2T_{\text{span}}}$. Note that the number of measurements is bounded by b_{yn} .

	M1	M2	M3	M4	M5
30 Bus	Y	N	N	N	Y
57 Bus	Y	N	N	N	Y
118 Bus	Y	N	N	N	Y
300 Bus	Y	N	N	N	Y
1354 Bus	Y	N	N	N	Y
2848 Bus	Y	N	N	N	Y

dependent. In contrast, due to the proposed linear basis, it is possible to provide a numerically-verifiable condition for checking whether the underlying set of clean measurements are enough to identify the actual state of the system. A spanning tree of a graph can be found in linear time by either depth-first search or breadth-first search. For instance, Fig. 2 illustrates a spanning tree of the IEEE 14 bus system. Moreover, it is a strong result in the sense that it does not depend on the numerical value of the true state (i.e., universally applicable).

Note that “identifiability” is a stronger condition than “observability.” The latter is usually based on DC-approximation (see [55] and [2, Chap. 4]), where a system is observable as long as the measurement Jacobian is nonsingular. However, for AC SE, “observability” implies the existence of a method (with potentially exponential-time complexity) to infer the state uniquely, but polynomial-time methods (such as local search)

may have a number of spurious local minima [13], [24]); however, “identifiability” indicates that the given method is guaranteed to efficiently recover the unique system state. In other words, “identifiability” is a sufficient condition for “observability.”

Remark 5. In general, there are $n_b - 1$ edges for a spanning tree with n_b nodes; hence, by the construction of the linear state can be recovered (either exactly when the dense noise representation in Sec. II-B, there are in total $n_b - 2$ number is negligible, or accurately enough for the case with dense of variables on the spanning tree. Therefore, as long as there are at least $3 - n_b$ independent measurements of branch flows or voltage magnitudes, we can achieve identifiability. This is empirically evaluated for several IEEE standard systems in Table I. It is well known that the traditional setting of power flow analysis with $2 - n_b$ measurements may have many spurious local minima [14], [30], [32]. However, it follows from the proposed linear basis that the network becomes identifiable with a limited but right set of measurements (e.g., M1 and M5 in Table I).

Remark 6 (Significance of Theorem 4) Local minima may not be a serious concern in normal situations where the state does not change rapidly and is within a normal operation region. However, the problem becomes crucial during adverse conditions where some states fall out of the normal operation region and need to be identified in real-time. In this case, if the initial point used by a local search algorithm is not close enough to the actual state, then the algorithm could arrive at a point that is a local minimum with no physical meaning. Mathematically speaking, SE can be regarded as an overdetermined power flow (PF) in the noiseless case. It is known that PF can have an exponential number of solutions for highly meshed networks [4], [14]. So, the SE in the noiseless case can have any number of solutions from 1 to a very large number. First, it is not known when the number of solutions is unique and that depends on how overdetermined the SE is. Second, although there was not much known about local solutions of optimal power flow (OPF) till 2009, there have been a lot of studies in the past 10 years showing how non-convex the problem is. The computational complexity of SE is an overlooked problem. The complexity of OPF is due to the nonlinearity of PF, and since PF is integral to SE, it could be highly nonconvex. Recent papers show that the problem could have many spurious solutions [30]. SE is a special case of the matrix sensing problem and the machine learning community has shown that the problem could be of high computational complexity with spurious solutions unless strong conditions are satisfied [33], [53], [54]. The focus of the power community has been mostly on how to treat noise and corrupted data in SE rather than dealing with its underlying non-convexity. Since the industry mostly uses DC models and various approximations for SE, they have not directly dealt with the complexity of the problem. Power community has also looked at SE/PF using nonlinear models and their concern has been on the convergence of local search algorithms. Since convergence could be to a wrong solution, the focus of the current work is to first understand when such wrong solutions never exist and second design an algorithm to find the correct solution.

Remark 7. Consider the measurement equation $y = Ax + b$, where $\text{supp}(b) = J$. Assume that the measurement matrix A satisfies the following conditions: (a) the lower eigenvalue is positive, i.e., $C_{\min}(J) > 0$; (b) the mutual incoherence condition $(J) < 1$ is satisfied. Then, the unique solution to (S1-L1), denoted as $(\hat{x}; \hat{b})$, is exact and recovers the true state (i.e., $\hat{x} = x$ and $\hat{b} = b$).

Remark 8. Consider the measurement equation $y = Ax + w + b$, where $\text{supp}(b) = J$ and w is a random vector with zero mean and subgaussian parameter ρ . Suppose that the rows of A are normalized, and that the measurement matrix satisfies the following conditions: (a) the lower eigenvalue is positive, (b) there exists a constant $\epsilon_0 > 0$ such that the mutual incoherence condition $(J) = 1 - \epsilon_0$. Let the regularization parameter λ be chosen such that

$$\lambda > \frac{2\rho}{n_m} \sqrt{2^2 \log n_m} \quad (9)$$

then, the following properties hold for the solution to (S1-LASSO) denoted as $(\hat{x}; \hat{b})$:

- 1) (No false inclusion) The solution $(\hat{x}; \hat{b})$ has no false bad data inclusion (i.e., $\text{supp}(\hat{b}) \subseteq \text{supp}(b)$) with probability greater than $1 - \frac{c_0}{n_m}$, for some constant $c_0 > 0$.
- 2) (Large bad data detection) Let

$$g(\lambda) = n_m \left(\frac{1}{2} \frac{\rho}{C_{\min}(J)} + \kappa_{\|b\|_b(Q_J^> Q_J)} \right) \|\hat{b}\|_b \kappa_1$$

be a threshold value. Then, all bad data measurements with magnitude greater than $g(\lambda)$ will be detected (i.e., if $\|b_j\|_b > g(\lambda)$, then $\|\hat{b}_j\|_b > 0$) with probability greater than $1 - \frac{c_1}{n_m}$ for some constant $c_1 > 0$.

- 3) (Bounded error) The estimator error is bounded by

$$\|\hat{x} - x\|_2 \leq \frac{\rho}{C_{\min}} \frac{\|x\|_2}{\lambda} + n_m \kappa_{\|x\|_x(Q_J^> Q_J)} \|\hat{b}\|_b \kappa_1 \kappa_2$$

with probability greater than $1 - \exp\left(-\frac{c_1 \lambda^2}{4}\right)$, where $\kappa_{\|x\|_x(Q_J^> Q_J)} \kappa_1 \kappa_2$ denotes $\|\cdot\|_2$ induced norm.

B. Global recovery conditions and error bounds

Despite the simplicity of the identifiability condition, real-world measurements are often subject to random sensor noise and sparse bad data whose support is often unknown. The

existence of dense noise) and estimation algorithms (i.e., (S1-L1) or (S1-LASSO)), it is remarkable that the global recovery conditions in Theorems 7 and 8 are coincident. In the case of negligible dense noise, then a strong global recovery is achieved, that is, both the true state and the bad data are detected. With the presence of dense noise, it is no longer possible to achieve exact recovery; however, Theorem 8 indicates that with a proper selection of the penalty coefficient λ , one can avoid false detection of bad data (part

is part 2), and achieve state estimation within bounded error margin. Furthermore, both the bad data threshold and the error

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bound decrease with stronger mutual incoherence condition and lower-eigenvalue condition.

The above analysis for (S1-LASSO) can be adapted to the case without dense noise, giving rise to the following corollary.

Corollary 9. Consider the measurement equation $y = Ax + b$, where $b \in \mathbb{R}^m$ has support J . Suppose that the rows of A are normalized, and the regularization parameter is chosen to be positive, i.e. $\lambda > 0$. Assume that A satisfies the following conditions: (a) the lower eigenvalue is positive, (b) the mutual incoherence condition $\mu(J) < 1$ is satisfied. Then, the following properties hold for the solution (S1-LASSO) denoted as $(\hat{x}; \hat{b})$:

- 1) (No false inclusion) The solution $(\hat{x}; \hat{b})$ has no bad data false inclusion (i.e. $\text{supp}(\hat{b}) \subseteq J$).
- 2) (Large bad data detection) Let $g(\lambda) = \frac{1}{n_m} \kappa_{\lambda} \|(Q_J^T Q_J)^{-1}\|_b^{-1} k_1$ be a threshold value. Then, all bad data measurements with magnitude greater than $g(\lambda)$ will be detected (i.e., if $b_{\lambda j} > g(\lambda)$, then $\hat{b}_j > 0$).
- 3) (Bounded error) The estimator error is bounded by

$$\|\hat{x} - x\|_2 \leq \frac{1}{\kappa_{\lambda}} \|(Q_J^T Q_J)^{-1}\|_b^{-1} k_1 \|\hat{b} - b\|_2$$

To understand the equivalence between Corollary 9 and Theorem 7, note that one can choose λ to be arbitrary close to 0 so the detection threshold and error bounds also approach 0. The proof of Theorem 8 is based on the primal-dual witness technique popularized by [52]. However, the key difference is that the existing literature in statistical learning only focuses on sparse signal recovery [49]–[52], while the present study needs to recover both the sparse signal (i.e., bad data) and the dense signal (i.e., regressor), which is technically more challenging to prove. Indeed, related works on this topic, such as robust principal component analysis [56] and dense error correction [57], employ different proof techniques than the present study.

In what follows, we will discuss the influence of the possible error in stage-1 estimation on the outcome of the second stage. Let the estimations \hat{x}^{re} and \hat{x}^{im} over a line $\mathbb{R}^2 \times L$ be given by:

$$\hat{x}^{re} = x^{re} + \epsilon^{re} \quad \text{and} \quad \hat{x}^{im} = x^{im} + \epsilon^{im};$$

where x^{re} and x^{im} are the true values, and ϵ^{re} and ϵ^{im} are the estimation errors from stage 1. We provide a bound on the phase estimation error for each bus $i \in N$.

Proposition 10. The estimation error of the phase $\hat{\theta}_k$ is bounded by the k -th component of the vector

$$(L^{-1} \mathbf{1})^T \mathbf{e};$$

where $\mathbf{e} \in \mathbb{R}^{n_x}$ is a vector with the elements $e_k = \frac{x^{re} \epsilon^{im} - x^{im} \epsilon^{re}}{x^{re} \hat{x}^{re}}$, and L is the matrix described in Sec. III-B.

To understand how the results of this section can guarantee the accurate recovery of the system state, we consider three different cases: (1) If there is no dense error or sparse bad data and the identifiability condition in Theorem 4 is satisfied, then it is guaranteed that there is a one-to-one correspondence between the proposed basis and the underlying voltage phasor.

(2) If there is no dense error but only sparse bad data, under the mutual incoherence condition in Theorems 7 and 8, it is guaranteed that the proposed basis can be uniquely recovered and it corresponds to the underlying voltage phasor. In the case with dense error, since it is theoretically impossible to exactly recover the underlying voltage, our goal is to obtain an estimate of the state that is as close to the ground truth as possible. In summary, the estimate is exactly equal to the actual state of the system in Cases 1 and 2, whereas there is a nonzero estimation error in Case 3 that is bounded in Proposition 10. It will be empirically shown that the proposed basis significantly outperforms the polar/rectangular basis by bypassing the non-convexity of the problem and finding a provably correct solution.

Due to the centrality of the mutual incoherence condition throughout the theoretical analysis, we provide an analysis on the likelihood of global recovery condition satisfaction using arguments based on concentration inequalities in probability.

C. Stochastic bound

In this analysis, we assume oblivious adversary, which indicates that the set is chosen uniformly at random. The goal is to gauge the likelihood that a random matrix with an arbitrary sparsity pattern will satisfy the mutual incoherence condition. It is important to capture the network-topology-induced pattern in A in this analysis.

Definition 11 (Sparsity pattern) For an arbitrary matrix $A \in \mathbb{R}^{n_m \times n_x}$, the sparsity pattern is a binary matrix $S \in \mathbb{R}^{n_m \times n_x}$ whose $(i; j)$ -th entry is equal to 1 only if $A_{ij} \neq 0$. Define the set of matrices with a given sparsity pattern as

$$\mathcal{S}(N) := \{A \in \mathbb{R}^{n_m \times n_x} \mid A \circ N = A\};$$

where \circ denotes the Hadamard (element-wise) product.

To conduct the analysis, we fix the sparsity pattern N and assume that A is a sparse matrix with the given pattern, where each entry is a random sub-Gaussian variable. In other words, $A = N \circ \mathbf{g}$; where $\mathbf{g} = [g_{ij}]_{i \in [n_m], j \in [n_x]}$ is a dense random matrix with independent and identically distributed sub-Gaussian random variables with variance proxy σ^2 (c.f., [58, Chap. 1] for a detailed account of the terminologies).

In the following, we introduce some metrics \mathcal{P} to measure the sparsity. For each $j \in [n_x]$, let $n_j^j = \sum_{i \in [n_m]} N_{ij}$ and $n_j^c = \sum_{i \in [n_m]} N_{ij}^c$ denote the numbers of nonzero entries in the columns of N_j and N_j^c , respectively, and let $n_j = \max_{i \in [n_m]} n_j^i$ and $n_j^c = \max_{i \in [n_m]} n_j^{i,c}$ be their upper bounds, respectively.

Theorem 12. Suppose that the following conditions hold:

- 1) (Bounded moments) There exist constants α_1, α_2 and β_1 such that $\mathbb{E} |g_{ij}|^q \leq \alpha_1 q^{\alpha_2}$ for every $i \in [n_m]$ and $j \in [n_x]$;
- 2) (Tall matrix) $n_m = (1 + \epsilon)n_x$ for some $\epsilon > 0$; and
- 3) (Saturated columns) There exists a constant $\beta_2 (0 < \beta_2 < 1)$ such that $n_j^c \leq \beta_2 \frac{n_j^j}{j^2}$ for all $i \in [n_m]$ and $j \in [n_x]$.

Then, the minimization (S1-L1) recovers the true state with changing rapidly. However, since both stages of the proposed probability (1) as long as the number of corrupted method are convex, the global optimal in each stage can be measurements j is not too high and satisfies the inequality found under mild conditions (Theorem 7 and Theorem 8).

$$\min c_4 j c_j; \frac{j c_j}{2 a_1 n_j n_x} \ln 2 n_x \ln \frac{2}{2 a_2}$$

where $a_1 = \frac{c^4}{32} (\frac{c^2}{64})^{\frac{1}{q-2}}$, $a_2 = \min f \frac{4}{c^2} a_1; 20 \frac{2}{g} \frac{\ln 2}{j c_j}$, and $\theta = \frac{c^2}{4 a_1}$.

The above theorem states that as long as the number of good measurements is greater than the number of non-zero elements in the spoiled part of the sensing matrix up to a constant multiplier (i.e., $j c_j \& \text{const } n_j n_x$), then with high probability, the mutual incoherence condition is satisfied. This agrees with the numerical results that the higher the number of measurements (with a fixed number of bad data), the smaller the mutual incoherence parameter.

D. Discussions

The premise of the proposed approach is that there are abundant data for the grid. However, we do not require that the data be highly reliable. In fact, with the prevalence of cheaply available sensors, we consider the important case that some of the data can be completely wrong due to various reasons such as sensor faults, communication errors or even cyber attacks. The situation of “abundant but untrusted” data is very different from the case of “redundant and reliable” data, because in the latter case, the data are believed to be high-fidelity with very low error rates. The former is also a more challenging scenario, and in particular classical SE based on nonlinear least squares fails due to the adverse influence of bad data. To resolve this issue, our algorithm is based on a novel set of linear basis, where the number of variables is larger than the number of real state variables; however, the minimum number of measurements to guarantee identifiability does not scale by the number of lines in the network, but by the number of buses (Theorem 4). This means that it can be a relatively succinct representation with the right set of measurements (i.e., available measurements can form a spanning tree of the network). Furthermore, our theoretical analysis shows that under some mild conditions, the algorithm recovers the ground truth even though part of the data are completely wrong (Theorems 7 and 8). Generally speaking, when the amount of data is low, a small set of bad data would deteriorate the estimation and even make the system unobservable. In other words, the prerequisite that the amount of measured data is relatively high is mainly for robustness purposes (Theorem 12).

Another advantage of the proposed method is the availability of a convex formulation that can be solved efficiently using second-order algorithms. The normal Gauss-Newton iteration for nonlinear least squares is nonconvex, and therefore it should be initialized at a point close to the ground truth to guarantee convergence. Since we do not know a priori the state of the system, there is no guarantee that the estimation will converge to a point close to the ground truth. This is especially true when there is an adverse situation with the state

Under the proposed framework, we do not explicitly distinguish leverage and/or vertical outliers [29], [59]. As long as the mutual incoherence condition is satisfied, the proposed method is able to deal with bad leverage points systematically. The mathematical framework of this paper is also different from two recent studies [39] and [60]. The work [39] focuses on an iterative algorithm by locally linearizing the nonlinear measurements; however, the convergence condition for nonlinear measurements needs to be verified using SDP and the analysis is based on the “almost Euclidean property” (similar to RIP), which is difficult to satisfy and the relation to the underlying model of the system is not clear. The work [60] also proposes a linear basis of representation, but it requires that current measurements be known for all lines; it also includes these measurements in the measurement Jacobian as if they are the true value, which can seriously bias the estimation if some of the measurements are bad data, also known as error-in-variables in the statistics literature [61]. By contrast, the sensing matrix A in our paper comes from system topology and physical parameters. In the normal situation, this matrix is close to the actual system; however, we also allow that some of the parameters in the matrix to be wrong, which means that the corresponding measurements are treated as bad data.

V. EXPERIMENTS

Numerical evaluations are performed on IEEE benchmark systems from MATPOWER [62]. This includes the Pan European Grid Advanced Simulation and State Estimation (PE-GASE) 9241-bus and 13659-bus systems, which represent the size and complexity of the European high voltage transmission network [63]. While PMU measurements can be incorporated in the proposed framework, unless otherwise stated, we assume the available measurements to include full nodal measurements (i.e., voltage magnitudes and real/reactive injections) and bidirectional real/reactive branch flows over all lines. All the experiments are performed on a personal laptop with 3.3GHz Intel Core i7 and 16GB memory.

In each case, we randomly generate 50 sets of dense noise and sparse bad data. The dense noise for each measurement is a zero-mean Gaussian variable, with standard deviation of 0.1 p.u. for voltage magnitude measurements and 0.1 p.u. for all the other measurements, where p.u. is the dense noise level. This setup is inspired by the fact that voltage magnitude sensors have higher standards of accuracy compared to power meters. For the sparse bad data, unless otherwise specified, its support is randomly selected among the line measurements, with the only assumption that at most one bad data measurement exists for each line. The values for the sparse noise can be arbitrarily large, and we assume that these parameters are uniformly chosen from the set $[4:25; 3:75][3:75; 4:25]$ (per unit).

We adopt the root-mean-square error (RMSE) as the performance metric, which is defined as $\frac{1}{n_b} \sum_{i=1}^{2N} |v_i - \hat{v}_i|^2$, where v_i and \hat{v}_i are the true and estimated complex voltage

(a) RMSE (b) lower eigenvalue λ_{\min} (left) and mutual incoherence (right)

Fig. 3: Evaluation of the (S1-L1)-cleaning-direct recovery (L1-Direct), (S1-LASSO)-cleaning-direct recovery (LASSO-Direct), and local search with ℓ_1 loss (L1-Local) and squared loss with cleaning step (since the RMSE is greater than 0.2, its line is not shown on the graph) for the IEEE 300-bus system. We vary the percentage of bad data measurements from 0% to 10% (out of all line measurements), with the dense noise level $\sigma_n = 0.5\%$. The plots in (b) indicate the median (line with circles) and the min/max value (shaded region).

TABLE II: Comparison of the (S1-L1)-cleaning-direct recovery (L1-Direct), (S1-LASSO)-cleaning-direct recovery (LASSO-Direct), and local search with ℓ_1 loss and Newton's method with bad data detection. We vary the percentage of bad data at 5% (out of all line measurements) and dense noise level $\sigma_n = 0.5\%$.

	Newton method			Local search			LASSO-Direct			L1-Direct		
	RMSE	F1	Time (s)	RMSE	F1	Time (s)	RMSE	F1	Time (s)	RMSE	F1	Time (s)
14 Bus	.002	.852	0.6	.001	1	0.3	.001	1	2.3	.001	1	2.2
30 Bus	.042	.808	2.4	.001	.996	0.4	.002	1	2.3	.002	1	2.2
57 Bus	.043	.827	3.2	.001	.998	1.2	.004	.999	2.3	.004	.999	2.1
118 Bus	.003	.848	7.4	.002	.980	4.1	.002	1	1.5	.002	1	1.3
300 Bus	.699	.379	58.1	.093	.858	21.6	.004	.999	2.6	.004	.999	1.2

at bus i . To evaluate the bad data detection accuracy of bad data percentage, the proposed methods can reliably we use the F1 score, which is defined as $\frac{2 \cdot \text{precision} \cdot \text{recall}}{\text{precision} + \text{recall}}$, recover the ground truth, even with 10% of arbitrarily bad data. This can be implied from the lower eigenvalue conditions and the mutual incoherence conditions, which remain well-conditioned with the presence of bad data. We also perform the experiments on other systems, as is shown in Table II with bad data fixed at 5% level and dense noise $\sigma_n = 0.5\%$. It can be observed that local search methods (with a cleaning step for Newton's method) perform relatively well when the scale is small (up to 118 buses), but the performance (e.g., RMSE (S1-L1) or (S1-LASSO) combined with stage-2 direct recovery and bad data detection F1 score) deteriorates significantly for method) with the current practice local search method using the squared loss Newton method, and another local search method that replaces the squared loss with ℓ_1 loss [30]. We use SeDuMi [64] as the linear programming solver, and the Next, we examine the performance of the proposed estimation MATLAB implementation of limited-memory BFGS [65] for estimators when both the dense noise and the bad data intensity vary. We test on the French very high voltage and high voltage transmission network with 2848 buses [63]. As is shown in Fig. 4, the algorithm achieves a low RMSE with up to 1000 bad data measurements and 1% level of dense noise. The detection score for bad data remains above 99% for all the scenarios. We also show that due to the high detection accuracy of the bad data, it is beneficial to redo the estimation after the cleaning stage (LASSO Clean), which can improve the RMSE of estimation especially when the number of bad data measurements is significant.

We compare the proposed method (stage-1 estimators (S1-L1) or (S1-LASSO) combined with stage-2 direct recovery and bad data detection F1 score) deteriorates significantly for method) with the current practice local search method using the squared loss Newton method, and another local search method that replaces the squared loss with ℓ_1 loss [30]. We use SeDuMi [64] as the linear programming solver, and the Next, we examine the performance of the proposed estimation MATLAB implementation of limited-memory BFGS [65] for estimators when both the dense noise and the bad data intensity vary. We test on the French very high voltage and high voltage transmission network with 2848 buses [63]. As is shown in Fig. 4, the algorithm achieves a low RMSE with up to 1000 bad data measurements and 1% level of dense noise. The detection score for bad data remains above 99% for all the scenarios. We also show that due to the high detection accuracy of the bad data, it is beneficial to redo the estimation after the cleaning stage (LASSO Clean), which can improve the RMSE of estimation especially when the number of bad data measurements is significant.

First, we evaluate the robustness of the methods to bad data. As is shown in Fig. 3, due to convergence issues and spurious local minima, none of the local search methods could correctly estimate the true state. On the other hand, with the increase in the largest system provided by MATPOWER and challenging

(a) RMSE for LASSO-Direct with (right) and without (left) cleaning step

(b) F1 score for (S1-LASSO)

Fig. 4: Evaluation of the (S1-LASSO)-direct recovery method on the PEGASE 2848-bus system. The dense noise level varies from 0 to 2%, and the number of bad data measurements ranges up to 2000 (roughly 9% of the total line measurements). The bad data detection accuracy is shown as the F1 score. After the detection of bad data, they are removed and the remaining clean data are used again in the estimation (LASSO Clean).

for the existing algorithms (Table III). We evaluate the performance of the algorithm with data redundancy lower than the full set of measurements. In particular, we first examine Case A where the measurements include voltage magnitudes at all buses v_k , $j, k \in \mathcal{N}$ and three branch flows over all lines $f_{ij}^{(c)}$; $p_{ij}^{(c)}$; $q_{ij}^{(c)}$, $g \in \mathcal{L}$ and then study Case B as a sparse scenario that includes voltage magnitudes at all buses v_k , $j, k \in \mathcal{N}$ but three branch flows on a spanning tree and n_b additional lines. The number of measurements scales by the number of buses in Case B, and is comparable to the classical setting of the power flow calculation. In this experiment, we relax the previous assumption of having only one bad measurement per line and instead allow all measurements associated with the line to be corrupted. This model allows accounting for correlated bad data, which are the most challenging ones to identify. The lines with compromised data are chosen randomly, but they are checked to ensure that the network does not become disconnected after removing those lines. We fix the dense noise level at 0.5% and the percentage of bad data at 1% of the full measurements (i.e., 0.01 $(3 n_b + 4 n_l)$ number of bad data) for all the cases. The number of bad data measurements ranges from 120 (for PEGASE 1354-bus) to 1228 (for PEGASE 13659-bus).

It can be observed that the performance is satisfactory in all of the scenarios, and the estimation becomes more robust with a higher number of measurements (i.e., lower RMSE and higher bad data detection accuracy). Moreover, all computations are performed within a minute, which is important for real-time situational awareness.

VI. CONCLUSION

In this study, we proposed a linear basis of representation for power system measurements that succinctly captures the topology of the network. This leads to a two-stage estimation approach that efficiently solve the nonconvex SE under mild conditions usually satisfied with a sufficient instrumentation of sensors. The proposed algorithm is provably robust to bad data. We developed a robustness metric based on a deterministic quantity called mutual incoherence. Theoretical analysis of

TABLE III: Evaluation in large-scale benchmarks. Case A includes voltage magnitude measurements on all buses and 3 branch flow measurements; Case B includes voltage measurement on all buses and 3 branch flow measurements on n_b lines that consist of a spanning tree of the network. We use the (S1-LASSO)-Cleaning-Direct recovery method in both cases.

	Case A		Case B		Time (sec)
	RMSE	F1	RMSE	F1	
1354 Bus	.003	.996	.003	.995	9.8
2848 Bus	.004	.995	.003	.996	13.6
3012 Bus	.003	.998	.001	.998	18.5
6495 Bus	.005	.994	.005	.996	36.2
9241 Bus	.007	.993	.009	.994	43.6
13659 Bus	.007	.994	.009	.995	52.4

the global recovery condition and statistical error bounds was conducted, which relied on this key metric. The algorithm demonstrated robustness to bad data in various empirical evaluations, and achieved superior performance compared to baselines. Above all, the proposed method exhibited a satisfactory scalability for large systems with more than 13,000 buses. In contrast to semidefinite programming relaxation approaches, the SE can be solved with high accuracy within a minute for such large systems. This can significantly improve real-time situational awareness of grid operation.

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APPENDIX

A. Proof of Theorem 4

Proof. Since the only element in the null space of $A_{\mathcal{J}}$ is 0, the measurement equation (4) has a unique solution, which corresponds to the true state x^* . By the algorithm in stage 2, outlined in Sec. III-B, this recovers the true state of the grid. \square

B. Proof of Theorem 7

Proof. The dual program of (S1-L1) is given by:

$$\max_{h \in \mathbb{R}^m} h^T y; \quad \text{subject to } h^T A = 0; \|h\|_1 = 1; \quad (L1\text{-Dual})$$

To show that $(x_\lambda; b_\lambda)$ is the optimal solution of (S1-L1), we simply need to find a dual certificate h_λ that satisfies the Karush-Kuhn-Tucker (KKT) conditions:

$$(\text{dual feasibility}) \quad h_\lambda^T A = 0; \quad (10)$$

$$(\text{stationarity}) \quad h_\lambda \in \partial \|b_\lambda\|_1; \quad (11)$$

where $\partial \|b_\lambda\|_1$ denotes the subgradient of $\|b_\lambda\|_1$. By the definition of $\mathcal{J} := \text{supp}(b_\lambda)$, we need to find h_λ such that $h_{\mathcal{J}} = \text{sign}(b_{\mathcal{J}})$ and $\|h_{\mathcal{J}^c}\|_1 = 1$. In fact, we can meet a slightly stronger condition for strict feasibility by choosing $h_{\mathcal{J}^c} = A_{\mathcal{J}^c}^+ A_{\mathcal{J}}^T \text{sign}(b_{\mathcal{J}})$, which satisfies

strict dual feasibility (i.e., $\|h_{\mathcal{J}^c}\|_1 < 1$) due to the mutual incoherence condition. Thus, this certifies the optimality of $(x_\lambda; b_\lambda)$ for (S1-L1).

To show that $(x_\lambda; b_\lambda)$ is the unique optimal solution, let $(x; b)$ be an arbitrary feasible point of (S1-L1) different from $(x_\lambda; b_\lambda)$. Due to the lower eigenvalue condition, the matrix $Q_{\mathcal{J}} := A_{\mathcal{J}}^+ A_{\mathcal{J}}^T$ has full column rank. Let $\mathcal{J} = \text{supp}(b)$, then \mathcal{J} must not be equal to or be a subset of \mathcal{J}_λ because otherwise, from $Q_{\mathcal{J}} x_{\mathcal{J}} = Q_{\mathcal{J}} x_{\mathcal{J}_\lambda} = y$, we must have

$$x_{\mathcal{J}} = x_{\mathcal{J}_\lambda}, \quad \text{which is contradictory to the assumption. Let}$$

$$\mathcal{J}^c = \mathcal{J} \setminus \mathcal{J}_\lambda, \quad \text{then,}$$

$$\|b_{\mathcal{J}^c}\|_1 < \|y\|_1 \quad (12)$$

$$= \|y\|_1 (A_{\mathcal{J}^c} x_{\mathcal{J}^c} + I_{\mathcal{J}^c}^T b_{\mathcal{J}^c}) \quad (13)$$

$$= \|y\|_1 \|b_{\mathcal{J}^c}\|_1 + \|y\|_1 \|b_{\mathcal{J}}\|_1 \quad (14)$$

$$\|y\|_1 \|b_{\mathcal{J}^c}\|_1 < \|y\|_1 \|b_{\mathcal{J}^c}\|_1 + \|y\|_1 \|b_{\mathcal{J}}\|_1 \quad (15)$$

$$< \|b_{\mathcal{J}^c}\|_1 + \|b_{\mathcal{J}}\|_1 \quad (16)$$

$$= \|b\|_1; \quad (17)$$

where (12) is due to the strong duality between (S1-L1) and (L1-Dual), (13) is due to the primal feasibility of $(x; b)$, (14) is due to the dual feasibility condition (10), (15) is due to the Hölder inequality, and (16) is due to the strict feasibility of $(x_\lambda; b_\lambda)$. Thus, we have shown the uniqueness of the optimal solution $(x_\lambda; b_\lambda)$. \square

C. Proof of Theorem 8

We design the primal-dual witness (PDW) process as follows (note that this is not an actual algorithm, because we do not know the true support \mathcal{J} ; rather, it is only part of a proof technique popularized by [52]):

- 1) Set $\hat{b}_{\mathcal{J}^c} = 0$
- 2) Determine $(x; \hat{b}_{\mathcal{J}})$ by solving the following program:

$$\min_{b \in \mathbb{R}^m; x \in \mathbb{R}^n} \frac{1}{2n_m} \|y - Ax - I_{\mathcal{J}}^T \hat{b}_{\mathcal{J}}\|_2^2 + \|b_{\mathcal{J}}\|_1; \quad (18)$$

and $\hat{z}_{\mathcal{J}} \in \partial \|b_{\mathcal{J}}\|_1$ satisfying

$$\frac{1}{n_m} I_{\mathcal{J}}^T (y - Ax - I_{\mathcal{J}}^T \hat{b}_{\mathcal{J}}) + \hat{z}_{\mathcal{J}} = 0; \quad (19)$$

$$A^T (y - Ax - I_{\mathcal{J}}^T \hat{b}_{\mathcal{J}}) = 0; \quad (20)$$

- 3) Solve $\hat{z}_{\mathcal{J}^c}$ via the zero-subgradient equation:

$$\frac{1}{n_m} (y - Ax - \hat{b}) + \hat{z} = 0 \quad (21)$$

and check whether strict feasibility condition $\|z_{\mathcal{J}^c}\|_1 < 1$ holds.

Lemma 13. If the PDW procedure succeeds, then $(x; \hat{b})$ where $\hat{b} = (\hat{b}_{\mathcal{J}}; 0)$ is unique optimal of (S1-LASSO).

Proof. If PDW succeeds, then the optimality conditions (20) and (21) are satisfied, which certify the optimality of $(x; \hat{b})$. The subgradient \hat{z} satisfies $\|z_{\mathcal{J}^c}\|_1 < 1$ and $\hat{z}; \hat{b} = \partial \|b\|_1$.

Now, let $(x; b)$ be any other optimal, and let $F(x; b) = \frac{1}{2n_m} \|y - Ax - bk\|_2^2$, then we have

$$F(x; \hat{b}) + \sum_{D \cup E} \lambda_j \hat{b}_j = F(x; b) + \sum_{D \cup E} \lambda_j b_j;$$

and hence we have

$$F(x; \hat{b}) + \sum_{D \cup E} \lambda_j \hat{b}_j - \sum_{D \cup E} \lambda_j b_j = F(x; b) + \sum_{D \cup E} \lambda_j b_j - \sum_{D \cup E} \lambda_j b_j;$$

By the optimality conditions (20) and (21), we have $\lambda_j = r_j F(x; \hat{b}) = \frac{1}{n_m} (y - A x - \hat{b})^T$ and $r_j F(x; \hat{b}) = 0$, which imply that

$$F(x; \hat{b}) - \sum_{D \cup E} \lambda_j F(x; \hat{b}) = \sum_{D \cup E} \lambda_j \hat{b}_j - \sum_{D \cup E} \lambda_j b_j = 0$$

due to convexity. We thus have $\lambda_j \hat{b}_j = \lambda_j b_j$. Since by

Holder's inequality, we also have $\lambda_j \hat{b}_j \leq k \lambda_j k_1$ and

$k \lambda_j k_1 \leq 1$, we must have $\lambda_j k_1 = \lambda_j b_j$, and $b_j = 0$ for $j \in J^c$.

This means that we have $\text{supp}(b) \subseteq \text{supp}(\hat{b}) \subseteq J$. By restricting the optimization \hat{b} in (S1-LASSO) to the support J and by the lower eigenvalue condition, the optimization is strictly convex and the uniqueness of the solution follows.

Lemma 14. Suppose that $Q_J^T Q_J$ is invertible for $J \subseteq [m]$, where $Q_J = [A \ I_J^T]$. Then, we have

$$(J) = k A_{J^c} I_x (Q_J^T Q_J)^{-1} I_b^T k_1; \quad (22)$$

Proof. We will show that for any given $J \subseteq [m]$, we have $A_{J^c} I_x (Q_J^T Q_J)^{-1} I_b^T = A_{J^c}^T A_J^T$. By the definition of Q_J and block matrix inversion formula, we have

$$\begin{aligned} & I_x (Q_J^T Q_J)^{-1} I_b^T \\ &= (A^T A)^{-1} A_J^T (I - A_J (A^T A)^{-1} A_J^T)^{-1} \\ &= (A^T A)^{-1} A_J^T (I + A_J (A_J^T A_J)^{-1} A_J^T) \\ &= (A^T A)^{-1} (I + A_J^T A_J (A_J^T A_J)^{-1}) A_J^T \\ &= (A_J^T A_J)^{-1} A_J^T; \end{aligned}$$

where the first equation follows from the Sherman–Morrison–Woodbury formula (c.f., Prop. 18 in Sec. G) and the rest are elementary operations. \square

Lemma 15. Suppose that $Q_J^T Q_J$ is invertible for a given $J \subseteq [m]$, where $Q_J = [A \ I_J^T]$. Then, we have

$$I_b (Q_J^T Q_J)^{-1} I_b^T = I + A_J (A_J^T A_J)^{-1} A_J^T \quad (23)$$

Proof. By the definition of Q_J and block matrix inversion formula, we have

$$\begin{aligned} I_b (Q_J^T Q_J)^{-1} I_b^T &= (I - A_J (A^T A)^{-1} A_J^T)^{-1} \\ &= I + A_J (A_J^T A_J)^{-1} A_J^T; \end{aligned}$$

where the second equation follows from the Sherman–Morrison–Woodbury formula. \square

Proof of Theorem 8

We prove each part sequentially:

Part 1): By the construction of PDW, we have $\lambda_{J^c} = b_{J^c} = 0$. The zero-subgradient condition (21) can be written as:

$$\begin{aligned} & \frac{1}{n_m} I_{J^c} A^T (x \setminus \hat{x}) + I_{J^c} (b \setminus \hat{b}) \\ & \frac{1}{n_m} I_{J^c} w \setminus + \lambda_{J^c} = 0; \end{aligned}$$

where the equations indexed by J^c can be re-written as:

$$\begin{aligned} & \frac{1}{n_m} I_{J^c} A^T I_J \hat{x} - I_{J^c} I_J^T x \setminus \hat{x} \\ & \frac{1}{n_m} I_{J^c} w \setminus + \lambda_{J^c} = 0; \end{aligned} \quad (24)$$

Solving for λ_{J^c} yields

$$\lambda_{J^c} = \frac{1}{n_m} I_{J^c} (A^T (x \setminus \hat{x}) + w \setminus); \quad (25)$$

Similarly, combining (20) and (24), we have

$$\begin{aligned} & \frac{1}{n_m} A^T A^T I_J \hat{x} - I_{J^c} I_J^T x \setminus \hat{x} \\ & \frac{1}{n_m} I_{J^c} w \setminus + \lambda_{J^c} = 0; \end{aligned}$$

Thus, by the lower eigenvalue condition (see Def. 1), we can

solve for the estimation error $\begin{bmatrix} x \setminus \hat{x} \\ b_{J^c} \setminus \hat{b}_{J^c} \end{bmatrix}$:

$$\begin{bmatrix} x \setminus \hat{x} \\ b_{J^c} \setminus \hat{b}_{J^c} \end{bmatrix} = (Q_J^T Q_J)^{-1} Q_J^T w \setminus + n_m (Q_J^T Q_J)^{-1} \begin{bmatrix} 0 \\ \lambda_{J^c} \end{bmatrix}; \quad (26)$$

Recall that I_x and I_b denote the matrix that consists of the first n_x rows and last j rows of the identity matrix of size $n_x + j$, respectively. Then, we can plug the estimation error of \hat{x} in (25) to get:

$$\begin{aligned} \lambda_{J^c} &= \frac{1}{n_m} I_{J^c} A^T I_x (Q_J^T Q_J)^{-1} Q_J^T w \setminus \\ &+ I_{J^c} A^T I_x (Q_J^T Q_J)^{-1} \begin{bmatrix} 0 \\ \lambda_{J^c} \end{bmatrix} + \frac{1}{n_m} I_{J^c} w \setminus \\ &= \underbrace{I_{J^c} A^T I_x (Q_J^T Q_J)^{-1} I_b^T}_{\{z\}} \lambda_{J^c} \\ &+ \underbrace{I_{J^c} I_x A^T I_x (Q_J^T Q_J)^{-1} Q_J^T}_{\{z\}} \frac{w \setminus}{n_m}; \end{aligned}$$

By the mutual incoherence condition (i.e., $\mu(J) = 1$ for $J \subseteq [m]$) and Lemma 14, we have $k_1 \leq 1$. Let $Q_J^T = Q_J^T Q_J (Q_J^T Q_J)^{-1} Q_J^T$ be the orthogonal projection matrix, then we have:

$$\begin{aligned} \lambda_{J^c} &= I_{J^c} Q_J^T + I_{J^c} I_J^T I_b (Q_J^T Q_J)^{-1} Q_J^T \frac{w \setminus}{n_m} \\ &= I_{J^c} Q_J^T + \frac{w \setminus}{n_m}; \end{aligned}$$

since $I_{J^c} I_J^T = 0$. Since the elements of $w \setminus$ are zero-mean

sub-Gaussian with parameter², and the projection operator Since by Prop. 17, has spectral norm one, we have

$$P \|k_{J^c} k_1\|_2 \leq 2j \exp\left(-\frac{n_m^2 t^2}{2}\right) :$$

Setting $t = \frac{1}{2}$ yields:

$$P \|k_{J^c} k_1\|_2 \leq 2 \exp\left(-\frac{n_m^2}{8} + \log(n_m + jJ)\right) :$$

By the design of Φ , we conclude that

$$P \|k_{J^c} k_1\|_2 \leq \frac{1}{2} \exp(-c_1 n_m^2) :$$

Part 2): Now, we will bound the estimation error in (26). First, we bound the infinity norm of $\hat{b}_J - b_J = l_b$. By triangle inequality,

$$\|l_b\|_\infty \leq \|l_b(Q_J^T Q_J)^{-1} Q_J^T w \setminus k_1\|_\infty + n_m \|l_b(Q_J^T Q_J)^{-1}\|_\infty k_1 :$$

Since the second term is deterministic, we will now bound the first term. By the normalized measurement condition (6) and the lower eigenvalue condition (5), each entry of $(Q_J^T Q_J)^{-1} Q_J^T w \setminus k_1$ is zero-mean sub-Gaussian with parameter at most

$$\sqrt{2} \sqrt{k(Q_J^T Q_J)^{-1} k_2} \frac{1}{C_{\min}}$$

Thus, by the union bound, we have

$$P \|l_b(Q_J^T Q_J)^{-1} Q_J^T w \setminus k_1\|_\infty > t \leq 2 \exp\left(-\frac{C_{\min} t^2}{2} + \log jJ\right) :$$

Then, setting $t = \frac{\beta_m}{2 C_{\min}}$, and note that by our choice of β_m , we have $\frac{C_{\min} t^2}{2} > \log jJ$. Thus, we conclude that

$$\|k_{J^c} \hat{b}_J - k_{J^c} b_J\|_\infty \leq n_m \frac{1}{2 C_{\min}} + \|l_b(Q_J^T Q_J)^{-1}\|_\infty k_1$$

with probability greater than $1 - 2 \exp(-c_2 n_m^2)$. This indicates that for bad data entries greater than

$$g(\cdot) = n_m \frac{1}{2 C_{\min}} + \|l_b(Q_J^T Q_J)^{-1}\|_\infty k_1$$

will be detected by \hat{b}_J .

Part 3): Now, we bound the ℓ_2 norm of the signal error $x \setminus k = l_x$,

$$\|l_x\|_2 \leq \sqrt{2} \sqrt{k(Q_J^T Q_J)^{-1} Q_J^T w \setminus k_2} + n_m \|l_x(Q_J^T Q_J)^{-1}\|_\infty k_1 :$$

For the first term, by the application of standard sub-Gaussian concentration (see Theorem 16), we have

$$P \|l_x(Q_J^T Q_J)^{-1} Q_J^T w \setminus k_2\|_2 > \sqrt{2} \sqrt{k(Q_J^T Q_J)^{-1} Q_J^T k_F} + t \|l_x(Q_J^T Q_J)^{-1} Q_J^T k_2\|_2 \exp\left(-\frac{C_1 t^2}{4}\right) :$$

$$\|l_x(Q_J^T Q_J)^{-1} Q_J^T k_F\|_2 \leq \sqrt{\frac{2}{n_m + jJ}} \sqrt{k(Q_J^T Q_J)^{-1} k_2} \sqrt{k(Q_J^T Q_J)^{-1} k_F} \sqrt{C_{\min}}$$

due to the lower eigenvalue condition (5) and the normalized measurement assumption (6), and similarly we have

$$\|l_x(Q_J^T Q_J)^{-1} Q_J^T k_2\|_2 \leq \sqrt{\frac{2}{n_m + jJ}} \sqrt{k(Q_J^T Q_J)^{-1} k_2} \sqrt{k(Q_J^T Q_J)^{-1} k_F} \sqrt{C_{\min}}$$

we have

$$P \|l_x(Q_J^T Q_J)^{-1} Q_J^T w \setminus k_2\|_2 > t \frac{1}{C_{\min}} \exp\left(-\frac{C_1 t^2}{4}\right) :$$

Together, we conclude that

$$P \|k \setminus k_2\|_2 \leq t \frac{1}{C_{\min}} \sqrt{\frac{2}{n_m + jJ}} + n_m \|l_x(Q_J^T Q_J)^{-1}\|_\infty k_1 \leq \frac{1}{2} \exp\left(-\frac{C_1 t^2}{4}\right)$$

D. Proof of Corollary 9

The proof is similar to that of Theorem 8. We need to make changes such that $w \setminus = 0$ whenever necessary, and some elementary operations lead to the results.

E. Proof of Proposition 10

Proof. The j -th component of the vector $\hat{\alpha}^j$:

$$[\hat{\alpha}^j]_j = \arctan \frac{x^{j \text{im}}}{x^{j \text{re}}} + \frac{x^{j \text{im}} x^{j \text{re}}}{x^{j \text{re}} x^{j \text{re}}}$$

Since the arctangent is a Lipschitz function with constant $\frac{1}{2}$, we can establish the bound:

$$|j[\hat{\alpha}^j]_j - [j]_j| \leq \frac{1}{2} \frac{x^{j \text{im}} x^{j \text{re}}}{x^{j \text{re}} x^{j \text{re}}} = |j| \epsilon$$

After using the closed-form expression (7) for the result will easily follow. \square

F. Proof of Theorem 12

Proof. For a vector $r \in \mathbb{R}^{2f+1}$; $1 \leq j \leq f$: From the definition of sub-Gaussian distribution, we have that for any $j \in [n_x]$:

$$E \exp(t \sum_{k=1}^j r_k) = \prod_{k=1}^j E \exp(t r_k) = \prod_{k=1}^j \exp\left(\frac{t^2}{2}\right)$$

Therefore, due to the symmetry of

$$r_j \sim \text{subG}(n_j^{-1/2});$$

and r is a sub-Gaussian random vector with variance proxy n_j^{-2} :

It is sufficient to have $k_A^j \geq r k_1 \geq 1$ to guarantee the perfect recovery. We further relax this condition to the form

