Global and Local Analyses of Nonlinear Low-Rank Matrix Recovery Problems

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Abstract

The restricted isometry property (RIP) is a well-known condition that guarantees the absence of spurious local minima in low-rank matrix recovery problems with linear measurements. In this paper, we introduce a novel property named bound difference property (BDP) to study low-rank matrix recovery problems with nonlinear measurements. Using RIP and BDP jointly, we first focus on the rank-1 matrix recovery problem, for which we propose a new criterion to certify the nonexistence of spurious local minima over the entire space. We then analyze the general case with an arbitrary rank and derive a condition to rule out the possibility of having a spurious solution in a ball around the true solution. The developed conditions lead to much stronger theoretical guarantees than the existing bounds on RIP.

1. Introduction

The low-rank matrix recovery problem plays a central role in many machine learning problems, such as recommendation systems (Koren et al., 2009) and motion detection (Zhou et al., 2013; Fattahi and Sojoudi, 2020). It also appears in engineering problems, such as power system state estimation (Zhang et al., 2018c). The goal of this problem is to recover an unknown low-rank matrix $M^* \in \mathbb{R}^{n \times n}$ from certain measurements of the entries of M^* . The measurement equations may be linear or nonlinear, which will be discussed separately in the next two subsections.

1.1 Low-Rank Matrix Recovery with Linear Measurements

The basic form of the low-rank matrix recovery problem is the symmetric and noiseless one with linear measurements and the quadratic loss. The linear measurements can be represented by a linear operator $\mathcal{A}: \mathbb{R}^{n \times n} \to \mathbb{R}^m$ given by

$$\mathcal{A}(M) = (\langle A_1, M \rangle, \dots, \langle A_m, M \rangle)^T.$$

^{*.} A preliminary version of this paper has appeared in Bi and Lavaei (2020). Compared with the conference paper, we have developed a major new result on the local guarantee for the absence of spurious local minima in the general rank-r case and included a new application from machine learning to illustrate the effectiveness of our results.

The ground-truth matrix M^* is assumed to be symmetric and positive semidefinite with $\operatorname{rank}(M^*) \leq r$. The recovery problem can be formulated as follows:

$$\min \quad \frac{1}{2} \|\mathcal{A}(M) - d\|^2$$
s.t.
$$\operatorname{rank}(M) \le r,$$

$$M \succeq 0, \quad M \in \mathbb{R}^{n \times n},$$

$$(1)$$

where $d = \mathcal{A}(M^*)$. By factoring the decision variable M into its low-rank factors XX^T , the above problem can be rewritten as the unconstrained problem:

$$\min_{X \in \mathbb{R}^{n \times r}} \left\{ \frac{1}{2} \| \mathcal{A}(XX^T) - d \|^2 \right\}.$$
(2)

The optimization (2) is commonly solved by local search methods. Since (2) is generally nonconvex, local search methods may converge to a spurious local minimum (a non-global local minimum is called a spurious solution). To provide theoretical guarantees on the performance of local search methods for the low-rank matrix recovery, several papers have developed various conditions under which the optimization (2) is free of spurious local minima. In what follows, we will briefly review the state-of-the-art results on this problem.

Given a linear operator \mathcal{A} , define its corresponding quadratic form $\mathcal{Q} : \mathbb{R}^{n \times n} \times \mathbb{R}^{n \times n} \to \mathbb{R}$ as

$$[\mathcal{Q}](K,L) = \langle \mathcal{A}(K), \mathcal{A}(L) \rangle, \tag{3}$$

for all $K, L \in \mathbb{R}^{n \times n}$.

Definition 1 (Recht et al. (2010)) A quadratic form $Q : \mathbb{R}^{n \times n} \times \mathbb{R}^{n \times n} \to \mathbb{R}$ satisfies the restricted isometry property *(RIP)* of rank 2r for a constant $\delta \in [0, 1)$, denoted as δ -RIP_{2r}, if

$$(1-\delta) \|K\|_F^2 \le [\mathcal{Q}](K,K) \le (1+\delta) \|K\|_F^2$$

for all matrices $K \in \mathbb{R}^{n \times n}$ with $\operatorname{rank}(K) \leq 2r$.

Ge et al. (2017) showed that the problem (2) has no spurious local minima if the quadratic form Q satisfies δ -RIP_{2r} with $\delta < 1/5$. Zhang et al. (2019) strengthened this result for the special case of r = 1 by showing that δ -RIP_{2r} with $\delta < 1/2$ is sufficient to guarantee the absence of spurious local minima for (2). Zhang et al. (2018a) provided an example with a spurious local minimum in case of $\delta = 1/2$ to support the tightness of the bound.

When $\delta \geq 1/5$ in the case r > 1 or $\delta \geq 1/2$ in the case r = 1, the δ -RIP_{2r} property can still be useful in the sense that they can lead to *local* guarantees for the absence of spurious local minima in the problem (2), rather than global guarantees. This means that there is no spurious local minimizer X as long as XX^T is in a neighborhood of M^* . Under a local guarantee, local search methods would still converge to the global optimal solution if they are initialized sufficiently close to the ground truth. Many techniques for finding such a good initial point have been developed in the literature (see Section 2 for a brief discussion). In the rank-1 case, Zhang et al. (2019) proved the following local guarantee for the absence of spurious local minima: **Theorem 2 (Zhang et al. (2019))** Assume that r = 1 and the quadratic form Q satisfies the δ -RIP₂ property for some constant δ such that

$$\delta < \sqrt{1 - \frac{\epsilon^2}{2(1-\epsilon)}}$$

with $0 < \epsilon \leq (\sqrt{5}-1)/2$. Then, the problem (2) has no spurious local minimizer X that satisfies

$$||XX^T - M^*||_F \le \epsilon ||M^*||_F.$$

Note that $\epsilon \|M^*\|_F$ defines the radius of the ball around the ground truth that is devoid of spurious solutions. The recent work (Zhang and Zhang, 2020) generalized the techniques in Zhang et al. (2019), which led to the following result that can be applied to the rankr case for any r but is weaker than Theorem 2 in the rank-1 case (rewritten here in an equivalent form):

Theorem 3 (Zhang and Zhang (2020)) Assume that the quadratic form Q satisfies the δ -RIP_{2r} property for some constant δ such that $\delta < \sqrt{1-\epsilon}$ with $0 < \epsilon \leq 1$. Then, the problem (2) has no spurious local minimizer X that satisfies

$$||XX^T - M^*||_F \le \epsilon \lambda_r(M^*).$$

1.2 Nonlinear Low-Rank Matrix Recovery

Given the above-mentioned results for the low-rank matrix recovery problems with linear measurements, it is natural to investigate whether these results can be extended to problems that are similar to (2) but have more complex objective functions. The purpose of this paper is to study the existence of spurious local minima for the general low-rank matrix recovery problem

$$\min_{X \in \mathbb{R}^{n \times r}} f(XX^T),\tag{4}$$

where $f : \mathbb{R}^{n \times n} \to \mathbb{R}$ is an arbitrary function. Problems of this form are abound in many machine learning tasks (see Section 6 for an application). Moreover, every polynomial optimization problem can be formulated as such, and therefore the analysis of (4) enables the design of global optimization techniques for nonconvex polynomial optimization (Madani et al., 2017). In this paper, f is always assumed to be twice continuously differentiable. The problem (2) is a special case of (4) by choosing

$$f(M) = \frac{1}{2} \|\mathcal{A}(M) - d\|^2.$$
(5)

In the case with linear measurements, note that $f(M^*) = 0$ and therefore M^* is a global minimizer of f. In other words, there are often infinitely many minimizers for f, but the goal is to find the ground-truth low-rank solution M^* . Similar to the linear measurement case, we assume that the problem (4) has a ground truth $M^* = ZZ^T$ with rank $(M^*) \leq r$ that is a global minimizer of f(M). The Hessian of the function f in (4), denoted as $\nabla^2 f(M)$, can be also regarded as a quadratic form whose action on any two matrices $K, L \in \mathbb{R}^{n \times n}$ is given by

$$[\nabla^2 f(M)](K,L) = \sum_{i,j,k,l=1}^n \frac{\partial^2 f}{\partial M_{ij} \partial M_{kl}}(M) K_{ij} L_{kl}$$

If f is considered to be the special function in (5), then its corresponding Hessian $\nabla^2 f(M)$ becomes exactly the quadratic form Q defined in (3). Therefore, we naturally extend the definition of the δ -RIP_{2r} property for quadratic forms given in Definition 1 to general functions f by restricting their Hessian.

Definition 4 A twice continuously differentiable function $f : \mathbb{R}^{n \times n} \to \mathbb{R}$ satisfies the restricted isometry property of rank 2r for a constant $\delta \in [0, 1)$, denoted as δ -RIP_{2r}, if

$$(1-\delta)\|K\|_F^2 \le [\nabla^2 f(M)](K,K) \le (1+\delta)\|K\|_F^2 \tag{6}$$

for all matrices $M, K \in \mathbb{R}^{n \times n}$ with $\operatorname{rank}(M) \leq 2r$ and $\operatorname{rank}(K) \leq 2r$.

It is still unknown whether the δ -RIP_{2r} condition could lead to the nonexistence of spurious local minima. However, Li et al. (2019) proved that the problem (4) has no spurious local minima under a stronger condition, named δ -RIP_{2r,4r} with $\delta < 1/5$, as defined below.

Definition 5 A twice continuously differentiable function $f : \mathbb{R}^{n \times n} \to \mathbb{R}$ satisfies the restricted isometry property of rank (2r, 4r) for a constant $\delta \in [0, 1)$, denoted as δ -RIP_{2r,4r}, if

$$(1-\delta) \|K\|_F^2 \le [\nabla^2 f(M)](K,K) \le (1+\delta) \|K\|_F^2$$

for all matrices $M, K \in \mathbb{R}^{n \times n}$ with $\operatorname{rank}(M) \leq 2r$ and $\operatorname{rank}(K) \leq 4r$.

For the general recovery problem (4) with r = 1, the previous results in Zhang et al. (2019) and Li et al. (2019) both have serious limitations. The bound $\delta < 1/2$ given in Zhang et al. (2019) is proven to be tight in the case when f is generated by linear measurements, but it is not applicable to nonlinear measurements. The bound $\delta < 1/5$ given in Li et al. (2019) can be applied to a general function f, but it is not tight even in the linear case. To address these issues, we develop a new criterion to guarantee the absence of spurious local minima globally in (4) for a general function f in the rank-1 case, which is more powerful than the previous conditions. Unlike the bound given in Li et al. (2019), our new criterion completely depends on the properties of the Hessian of the function f applied to rank-2 matrices, rather than rank-4 matrices. Note that the rank-1 case has applications in many problems, such as motion detection (Fattahi and Sojoudi, 2020) and power system state estimation (Zhang et al., 2018c).

For the problem (4) with r > 1, we also present a new local guarantee for the absence of spurious local minima near the ground truth M^* . Our work not only offers the first result in the literature on the local guarantee for general nonlinear problems, but also exhibits an improvement over the previous results stated in Theorem 2 and Theorem 3 for problems with linear measurements.

1.3 Notations

 I_n is the identity matrix of size $n \times n$, and $\operatorname{diag}(a_1, \ldots, a_n)$ is a diagonal matrix whose diagonal entries are a_1, \ldots, a_n . $\mathbf{A} = \operatorname{vec} A$ is the vector obtained from stacking the columns of a matrix A. Given a vector $\mathbf{A} \in \mathbb{R}^{n^2}$, define its symmetric matricization $\operatorname{mat}_S \mathbf{A} = (A + A^T)/2$, where $A \in \mathbb{R}^{n \times n}$ is the unique matrix satisfying $\mathbf{A} = \operatorname{vec} A$. $A \otimes B$ is the Kronecker product of A and B, which satisfies the well-known identity:

$$\operatorname{vec}(AXB^T) = (B \otimes A) \operatorname{vec} X.$$

For two matrices A, B of the same size, define their inner product

$$\langle A, B \rangle = \operatorname{tr}(A^T B) = \langle \operatorname{vec} A, \operatorname{vec} B \rangle,$$

and let $||A||_F = \sqrt{\langle A, A \rangle}$ denote the Frobenius norm of the matrix A, where tr(·) is the trace operator. Moreover, ||v|| is the Euclidean norm of the vector v. For a square matrix $A \in \mathbb{R}^{n \times n}$, $A \succeq 0$ means that A is symmetric and positive semidefinite. Let

$$\lambda_1(A) \ge \lambda_2(A) \ge \dots \ge \lambda_n(A)$$

denote the eigenvalues of A sorted in a decreasing order.

2. Related Works

The classical approach for solving low-rank matrix recovery problems is through convex relaxations. A semidefinite program can be obtained by removing the nonconvex low-rank constraint in (1). Since the seminal work by Recht et al. (2010), there is a plethora of researches on deriving conditions under which the convex relaxation is able to recover the exact solution of the original nonconvex problem. Most of the proposed conditions are based on RIP, including Candès and Plan (2011); Candes et al. (2013); Cai and Zhang (2013); Zhang and Li (2017); Li et al. (2020). Another direction is to show that the convex relaxation is exact with high probability if the measurements are random and have sufficient number of samples (Candès and Recht, 2009; Candès and Tao, 2010). The major drawback of the convex relaxation approach is that semidefinite programs are expensive to solve for large-scale problems arising in machine learning.

An alternative approach for solving low-rank matrix recovery problems is based on local search methods, such as gradient descent algorithms (Rennie and Srebro, 2005; Lee et al., 2010; Recht and Ré, 2013; Ge et al., 2015; Tu et al., 2016), iterative hard thresholding (Rauhut et al., 2017) and trust-region methods (Sun et al., 2016; Boumal et al., 2019). They can be efficiently applied to large-scale problems, but the quality of the obtained solution depends on whether the objective has a benign landscape. As previously discussed in Section 1, RIP-type conditions can be used to guarantee the absence of spurious local minima over the entire space. Under these conditions, any local search method that converges to a local minimum will be able to recover the globally optimal solution. The existing proof techniques for the analysis of spurious local minima can be roughly categorized into two groups: 1) checking whether an arbitrary matrix X is a spurious local minimizer for a given measurement operator \mathcal{A} and a given ground truth matrix M^* (Bhojanapalli et al., 2016b; Ge et al., 2016, 2017; Park et al., 2017; Zhang et al., 2018b; Li et al., 2019); 2) checking whether a measurement operator \mathcal{A} exists that makes a given point X a spurious local solution for a given ground truth matrix M^* . The second proof technique was first proposed in Zhang et al. (2019) and later extended in Molybog et al. (2020); Zhang and Zhang (2020) for structured operators. We adopt the same technique in the current paper.

In the case when the absence of spurious local minima in the entire space cannot be guaranteed, there are various approaches to handle the problem: 1) apply special initialization schemes such as spectral methods to find an initial point sufficiently close to the ground truth (Zheng and Lafferty, 2015; Candes et al., 2015; Bhojanapalli et al., 2016a; Sun and Luo, 2016; Park et al., 2018); 2) use randomized algorithms such as stochastic descent methods to escape saddle points or poor local minimizers (Ge et al., 2015); or 3) initialize the algorithm randomly multiple times (Goldstein and Studer, 2018; Zhang and Zhang, 2020).

3. Main Results

To obtain a powerful condition for guaranteeing the absence of spurious local minima in problem (4), it is helpful to shed light on a distinguishing property of the function in (5) for linear measurements that does not hold in the general case: the Hessian matrices at all points are equal. If a general function f satisfies δ -RIP_{2r}, (6) intuitively states that the Hessian $\nabla^2 f(M)$ should be close to the quadratic form defined by an identity matrix, at least when applied to rank-2r matrices. Hence, $\nabla^2 f(M)$ should change slowly when Malters. The above discussion motivates the introduction of a new notion below.

Definition 6 A twice continuously differentiable function $f : \mathbb{R}^{n \times n} \to \mathbb{R}$ satisfies the bounded difference property of rank 2r for a constant $\kappa \geq 0$, denoted as κ -BDP_{2r}, if

$$\|[\nabla^2 f(M) - \nabla^2 f(M')](K, L)\| \le \kappa \|K\|_F \|L\|_F$$
(7)

for all matrices $M, M', K, L \in \mathbb{R}^{n \times n}$ whose ranks are at most 2r.

It turns out that the RIP and BDP properties are not fully independent. Their relationship is summarized in the following theorems that will be proved in Section 4.

Theorem 7 If the function f satisfies δ -RIP_{2r}, then it also satisfies 4δ -BDP_{2r}.

Theorem 8 If the function f satisfies δ -RIP_{2r.4r}, then it also satisfies 2δ -BDP_{2r}.

The bounds in the above two theorems are tight. In Section 4, we will construct a class of functions f that satisfy the δ -RIP_{2r} property but do not satisfy the κ -BDP_{2r} property for some κ with κ/δ being arbitrarily close to 4. Similar examples can also be constructed for Theorem 8.

The main results of this paper will be stated below, which are powerful criteria for the global and local nonexistence of spurious local minima. The proofs are given in Section 5.

Theorem 9 (Global Guarantee for r = 1) When r = 1, the problem (4) has no spurious local minima if the function f satisfies the δ -RIP₂ and κ -BDP₂ properties for some constants δ and κ such that

$$\delta < \frac{2 - 6(1 + \sqrt{2})\kappa}{4 + 6(1 + \sqrt{2})\kappa}.$$

Theorem 10 (Local Guarantee for $r \ge 1$) Assume that the function f satisfies the δ -RIP_{2r} property for some constant δ such that

$$\delta < \sqrt{1 - \frac{3 + 2\sqrt{2}}{4}\epsilon^2}$$

with $0 < \epsilon \leq 2(\sqrt{2}-1)$. Then, the problem (4) has no spurious local minimizer X that satisfies

$$||XX^T - M^*||_F \le \epsilon \lambda_r(M^*).$$

In the case of linear measurements and the quadratic loss, the function f satisfies the κ -BDP_{2r} property with $\kappa = 0$. Hence, Theorem 9 immediately recovers the result in Zhang et al. (2019) stating that the problem (2) with r = 1 has no spurious local minima if the quadratic form \mathcal{Q} satisfies the δ -RIP₂ property with $\delta < 1/2$.

As a by-product, Theorem 10 also improves the existing local guarantees summarized in Theorem 2 and Theorem 3 for certain linear cases. If r = 1, Theorem 10 can possibly offer a region free of spurious local minima that is larger than the region obtained from Theorem 2. The reason is that ϵ in Theorem 2 is capped at $(\sqrt{5}-1)/2$, which is increased to $2(\sqrt{2}-1)$ in Theorem 10 (note that $\lambda_r(M^*) = ||M^*||_F$ if r = 1). For an arbitrary rank, Theorem 10 strengthens the result of Theorem 3 in terms of the order of the bound as a function of ϵ , when ϵ is small.

Theorem 9 and Theorem 10 are even more powerful for functions f associated with nonlinear measurements. At the end of Section 4 and in Section 6, we will offer such examples for which the absence of spurious local minima can be certified by Theorem 9 or Theorem 10, while the existing conditions in the literature fail to work.

4. RIP and BDP Properties

In this section, the relationship among the RIP_{2r} , $\text{RIP}_{2r,4r}$ and BDP_{2r} properties of a given function f will be investigated. We will first prove Theorem 7 and Theorem 8, and then show that the bounds in these theorems are tight. The following lemma will be needed, which appears in Candès (2008); Bhojanapalli et al. (2016b); Li et al. (2019) under different notations. We include a short proof here for completeness.

Lemma 11 If a quadratic form Q satisfies δ -RIP_{2r}, then

 $|[\mathcal{Q}](K,L) - \langle K,L \rangle| \le \delta ||K||_F ||L||_F$

for all matrices $K, L \in \mathbb{R}^{n \times n}$ with $\operatorname{rank}(K) \leq r$, $\operatorname{rank}(L) \leq r$.

Proof Without loss of generality, assume that $||K||_F = ||L||_F = 1$. By the δ -RIP_{2r} property of \mathcal{Q} , we have

$$(1-\delta) \|K-L\|_F^2 \le [\mathcal{Q}](K-L,K-L) \le (1+\delta) \|K-L\|_F^2, (1-\delta) \|K+L\|_F^2 \le [\mathcal{Q}](K+L,K+L) \le (1+\delta) \|K+L\|_F^2.$$

Taking the difference between the above two inequalities, one can obtain

$$4[\mathcal{Q}](K,L) \le (1+\delta) \|K+L\|_F^2 - (1-\delta) \|K-L\|_F^2 = 4\delta + 4\langle K,L\rangle, -4[\mathcal{Q}](K,L) \le (1+\delta) \|K-L\|_F^2 - (1-\delta) \|K+L\|_F^2 = 4\delta - 4\langle K,L\rangle,$$

which proves the desired inequality.

Proof of Theorem 8 Let M and M' be two matrices of rank at most 2r. By the definition of δ -RIP_{2r,4r} of the function f, both $\nabla^2 f(M)$ and $\nabla^2 f(M')$ satisfy δ -RIP_{4r}. After the constant r in the statement of Lemma 11 is replaced by 2r, we obtain

$$|[\nabla^2 f(M)](K,L) - \langle K,L\rangle| \le \delta ||K||_F ||L||_F,$$
$$|[\nabla^2 f(M')](K,L) - \langle K,L\rangle| \le \delta ||K||_F ||L||_F,$$

for all matrices $K, L \in \mathbb{R}^{n \times n}$ of rank at most 2r, which leads to (7) for $\kappa = 2\delta$.

Proof of Theorem 7 We first prove that any quadratic form \mathcal{Q} with δ -RIP_{2r} satisfies

$$|[\mathcal{Q}](K,L) - \langle K,L \rangle| \le 2\delta ||K||_F ||L||_F, \tag{8}$$

for all matrices $K, L \in \mathbb{R}^{n \times n}$ of rank at most 2r. Let $K = UDV^T$ be the singular value decomposition of K. Write $D = D_1 + D_2$ in which D_1 and D_2 both have at most r nonzero entries, and let $K_1 = UD_1V^T$ and $K_2 = UD_2V^T$. Then, $K = K_1 + K_2$, where rank $(K_1) \leq r$, rank $(K_2) \leq r$ and $\langle K_1, K_2 \rangle = 0$. We decompose $L = L_1 + L_2$ similarly. By Lemma 11, it holds that

$$\begin{split} |[\mathcal{Q}](K,L) - \langle K,L \rangle| &\leq |[\mathcal{Q}](K_1,L_1) - \langle K_1,L_1 \rangle| + |[\mathcal{Q}](K_1,L_2) - \langle K_1,L_2 \rangle| \\ &+ |[\mathcal{Q}](K_2,L_1) - \langle K_2,L_1 \rangle| + |[\mathcal{Q}](K_2,L_2) - \langle K_2,L_2 \rangle| \\ &\leq \delta(\|K_1\|_F + \|K_2\|_F)(\|L_1\|_F + \|L_2\|_F) \\ &\leq 2\delta \sqrt{\|K_1\|_F^2 + \|K_2\|_F^2} \sqrt{\|L_1\|_F^2 + \|L_2\|_F^2} \\ &= 2\delta \|K\|_F \|L\|_F. \end{split}$$

The remaining proof is exactly the same as the proof of Theorem 8.

The inequality (8) is parallel to the square root lifting inequality (Cai et al., 2010) in the compressed sensing problem. Our result can be regarded as a generalization of that result to the low-rank matrix recovery problem.

In what follows, we will show that the bounds in Theorem 7 and Theorem 8 are tight. To this end, we will work on examples of function f with δ -RIP_{2r} or δ -RIP_{4r} for a small δ whose Hessian has a large variation across different points. Consider an integer $n \ge 4$ and an integer $r \ge 1$. Let

$$A_1 = \frac{1}{\sqrt{n}} \operatorname{diag}(a_1, \dots, a_n)$$

with $a_i \in \{-1, 1\}$ whose exact value will be determined later. One can extend A_1 to an orthonormal basis A_1, \ldots, A_{n^2} of the space $\mathbb{R}^{n \times n}$. Define a linear operator $\mathcal{A} : \mathbb{R}^{n \times n} \to \mathbb{R}^{n^2-1}$ by letting

$$\mathcal{A}(M) = (\langle A_2, M \rangle, \dots, \langle A_{n^2}, M \rangle).$$

Then, for every matrix $M \in \mathbb{R}^{n \times n}$, it holds that

$$\|\mathcal{A}(M)\|^{2} = \|M\|_{F}^{2} - (\langle A_{1}, M \rangle)^{2} \le \|M\|_{F}^{2}.$$

Now, assume that M is a matrix with rank $(M) \leq 2r$, and let $\sigma_1(M), \ldots, \sigma_{2r}(M)$ denote its 2r largest singular values. Observe that

$$|\langle A_1, M \rangle| \le \frac{1}{\sqrt{n}} \sum_{i=1}^n |M_{ii}| \le \frac{1}{\sqrt{n}} \sum_{i=1}^{2r} \sigma_i(M) = \sqrt{\frac{2r}{n}} \sqrt{\sum_{i=1}^{2r} \sigma_i^2(M)} = \sqrt{\frac{2r}{n}} ||M||_F,$$

which implies that

$$\|\mathcal{A}(M)\|^2 = \|M\|_F^2 - (\langle A_1, M \rangle)^2 \ge \left(1 - \frac{2r}{n}\right) \|M\|_F^2.$$

Define a scaled linear operator $\overline{\mathcal{A}}$ as

$$\bar{\mathcal{A}}(M) = \sqrt{\frac{n}{n-r}} \mathcal{A}(M), \quad \forall M \in \mathbb{R}^{n \times n}.$$

Thus, the relation

$$\left(1 - \frac{r}{n-r}\right) \|M\|_F^2 \le \|\bar{\mathcal{A}}(M)\|^2 \le \left(1 + \frac{r}{n-r}\right) \|M\|_F^2 \tag{9}$$

holds for all $M \in \mathbb{R}^{n \times n}$ with rank $(M) \leq 2r$.

After choosing $A_1 = (1/\sqrt{n})I_n$ in the above argument, let \mathcal{A} be the resulting linear operator and \mathcal{Q} be the quadratic form in (3) that corresponds to the scaled linear operator $\overline{\mathcal{A}}$. By the same argument, a similar linear operator \mathcal{A}' and the corresponding quadratic form \mathcal{Q}' can be obtained after choosing

$$A'_{1} = \frac{1}{\sqrt{n}} \operatorname{diag}(1, 1, -1, -1, 1, \dots, 1).$$
(10)

Now, we select $K = \text{diag}(1, 1, 0, 0, 0, \dots, 0)$ and $L = \text{diag}(0, 0, 1, 1, 0, \dots, 0)$. Then,

$$|[\mathcal{Q} - \mathcal{Q}'](K, L)| = \frac{n}{n-r} |\langle \mathcal{A}(K), \mathcal{A}(L) \rangle - \langle \mathcal{A}'(K), \mathcal{A}'(L) \rangle|$$

$$= \frac{n}{n-r} |-\langle A_1, K \rangle \langle A_1, L \rangle + \langle A'_1, K \rangle \langle A'_1, L \rangle|$$

$$= \frac{4}{n-r} ||K||_F ||L||_F.$$
 (11)

In the case r = 1, it follows from (9) that both of the constructed quadratic forms Q and Q' satisfy δ -RIP₂ with $\delta = 1/(n-1)$. If one can find a twice continuously differentiable function f satisfying δ -RIP₂ such that

$$\nabla^2 f(M) = \mathcal{Q}, \quad \nabla^2 f(M') = \mathcal{Q}'$$

hold at two particular points $M, M' \in \mathbb{R}^{n \times n}$ with $\operatorname{rank}(M) \leq 2$ and $\operatorname{rank}(M') \leq 2$, then by (11) the function f cannot satisfy κ -BDP₂ for $\kappa < 4\delta$. Since the design of such function is cumbersome, we will use a weaker result that serves the same purpose. This result, to be formalized in Lemma 12, states that for every $\mu > 0$, one can find a twice continuously differentiable function f with $(\delta + \mu)$ -RIP₂ and two matrices $M, M' \in \mathbb{R}^{n \times n}$ of rank at most 1 satisfying the following inequalities:

$$|[\nabla^2 f(M) - Q](K, L)| \le \mu ||K||_F ||L||_F, [\nabla^2 f(M') - Q'](K, L)| \le \mu ||K||_F ||L||_F.$$
(12)

Combining (11) and (12) yields that

$$|[\nabla^2 f(M) - \nabla^2 f(M')](K, L)| \le (4\delta + 2\mu) ||K||_F ||L||_F$$

Therefore, the function f cannot satisfy the κ -BDP₂ property for any $\kappa < 4\delta + 2\mu$. Since μ can be made arbitrarily small, this shows that the constant 4δ in Theorem 7 cannot be improved. Similarly, by choosing r = 2 instead of r = 1 and repeating the above argument, one can show that the constant 2δ in Theorem 8 cannot be improved either.

Lemma 12 Consider two quadratic forms \mathcal{Q} and \mathcal{Q}' satisfying the δ -RIP_{2r} property. For every $\mu > 0$, there exists a twice continuously differentiable function $f : \mathbb{R}^{n \times n} \to \mathbb{R}$ and two matrices $M, M' \in \mathbb{R}^{n \times n}$ with rank $(M) \leq 1$ and rank $(M') \leq 1$ such that f satisfies the $(\delta + \mu)$ -RIP_{2r} property and that (12) holds for all $K, L \in \mathbb{R}^{n \times n}$.

Proof Given $\mu > 0$, let f be given as

$$f(V) = \frac{1}{2}[\mathcal{Q}'](V,V) + \frac{1}{2}H(||V||_F^2)[\Delta](V,V),$$

where $\Delta = \mathcal{Q} - \mathcal{Q}'$ and $H : \mathbb{R} \to \mathbb{R}$ is defined as

$$H(t) = \begin{cases} 0, & \text{if } t \le 0, \\ \exp(-1/t^{\gamma}), & \text{if } t > 0. \end{cases}$$

Here, $\gamma \in (0, 1)$ is a constant that will be determined later. It is straightforward to verify that H is twice continuously differentiable and

$$H'(0) = H''(0) = 0, (13a)$$

$$|tH'(t)| \le \frac{\gamma}{e}, \quad |t^2H''(t)| \le \frac{4\gamma}{e}, \quad \forall t \in \mathbb{R}.$$
 (13b)

The basic idea behind the above construction of f is that when γ is chosen to be small, the growth of the function H becomes so slow that it can be regarded as a constant when computing the Hessian of the above function f. As a result, the Hessian is approximately a linear combination of two quadratic forms Q and Q' with the δ -RIP_{2r} property. Formally, the Hessian $\nabla^2 f(V)$ of f at a particular matrix $V \in \mathbb{R}^{n \times n}$, when applied to arbitrary $K, L \in \mathbb{R}^{n \times n}$, is given by

$$[\nabla^{2} f(V)](K,L) = 2H''(\|V\|_{F}^{2})[\Delta](V,V)\langle V,K\rangle\langle V,L\rangle + H'(\|V\|_{F}^{2})[\Delta](V,V)\langle K,L\rangle + 2H'(\|V\|_{F}^{2})([\Delta](L,V)\langle V,K\rangle + [\Delta](K,V)\langle V,L\rangle) + [\mathcal{Q}' + H(\|V\|_{F}^{2})\Delta](K,L).$$

$$(14)$$

By compactness, there exists a constant C > 0 such that

$$|[\Delta](A,B)| \le C ||A||_F ||B||_F \tag{15}$$

holds for all $A, B \in \mathbb{R}^{n \times n}$. We choose a sufficiently small γ such that $26\gamma C/e \leq \mu$. By (13b), (14), (15) and the Cauchy-Schwartz inequality, we have

$$|[\nabla^2 f(V) - \mathcal{Q}' - H(\|V\|_F^2)\Delta](K,L)| \le \frac{13\gamma C}{e} \|K\|_F \|L\|_F \le \frac{\mu}{2} \|K\|_F \|L\|_F.$$
(16)

To prove that the function f satisfies $(\delta + \mu)$ -RIP_{2r}, assume for now that K = L and rank $(K) \leq 2r$. The inequality $0 \leq H(||V||_F^2) \leq 1$ and the δ -RIP_{2r} property of \mathcal{Q} and \mathcal{Q}' imply that

$$(1-\delta) \|K\|_F^2 \le [\mathcal{Q}' + H(\|V\|_F^2)\Delta](K,K) \le (1+\delta) \|K\|_F^2.$$

By (16) and the above inequality, the function f satisfies the $(\delta + \mu)$ -RIP_{2r} property. To prove the existence of M and M' satisfying (12), we select M' = 0 and

$$M = \operatorname{diag}(s, 0, \dots, 0).$$

For any $K, L \in \mathbb{R}^{n \times n}$, it follows from (13a) and (14) that

$$[\nabla^2 f(M') - Q'](K, L) = 0.$$
(17)

Moreover, (15) and (16) yield that

$$\begin{aligned} |[\nabla^2 f(M) - \mathcal{Q}](K, L)| &\leq \frac{\mu}{2} ||K||_F ||L||_F + |[\mathcal{Q}' + H(||M||_F^2)\Delta - \mathcal{Q}](K, L)| \\ &\leq \left(\frac{\mu}{2} + (1 - H(||M||_F^2))C\right) ||K||_F ||L||_F. \end{aligned}$$

Since $H(||M||_F^2) \to 1$ as $s \to +\infty$, (12) is satisfied as long as s is sufficiently large.

The above argument also provides examples of the function f whose corresponding recovery problem (4) can be certified to have no spurious local minima via Theorem 9, while the existing results in the literature fail to do so. Following the above construction, choose n = 4, r = 1, and let

$$\tilde{f}(V) = \frac{1-\lambda}{2} [\mathcal{Q}'](V, V) + \lambda f(V),$$

for some $\lambda \in [0, 1]$. The Hessian can be written as

$$\nabla^2 \tilde{f}(V) = (1 - \lambda)\mathcal{Q}' + \lambda \nabla^2 f(V).$$
(18)

If $\lambda > 0$, the Hessian of \tilde{f} is not a constant, and therefore the condition in Zhang et al. (2019) cannot be applied. On the other hand, it follows from (17) that

$$[\nabla^2 \tilde{f}(0)](A'_1, A'_1) = [\mathcal{Q}'](A'_1, A'_1) = 0,$$

for the matrix A'_1 of rank 4 defined in (10). Thus, the function \tilde{f} cannot satisfy the δ -RIP_{2,4} property for any $\delta \in [0, 1)$. This implies that the condition in Li et al. (2019) cannot be applied either. In contrast, note that the quadratic form \mathcal{Q}' satisfies the 1/3-RIP₂ property and the function f satisfies the $(1/3 + \mu)$ -RIP₂ property. Therefore, it can be concluded from (18) that the function \tilde{f} also satisfies the $(1/3 + \mu)$ -RIP₂ property. In light of Theorem 7, f satisfies $4(1/3 + \mu)$ -BDP₂ and thus f' satisfies $4\lambda(1/3 + \mu)$ -BDP₂. Hence, Theorem 9 certifies the absence of spurious local minima as long as λ and μ jointly satisfy

$$\frac{1}{3} + \mu < \frac{2 - 6(1 + \sqrt{2})4\lambda(1/3 + \mu)}{4 + 6(1 + \sqrt{2})4\lambda(1/3 + \mu)}.$$

5. Proofs of Main Results

Our approach consists of two major steps. The first step is to find necessary conditions that the function f with the δ -RIP_{2r} and κ -BDP_{2r} properties must satisfy if the corresponding problem (4) has a local minimizer X such that $XX^T \neq M^*$, where M^* is the ground truth. The second step is to develop certain conditions on δ and κ that rule out the satisfaction of the above necessary condition.

Before proceeding with the proofs, we need to introduce some notations. Given two matrices $X, Z \in \mathbb{R}^{n \times r}$, define

$$\mathbf{e} = \operatorname{vec}(XX^T - ZZ^T) \in \mathbb{R}^{n^2},$$

and let $\mathbf{X} \in \mathbb{R}^{n^2 \times nr}$ be the matrix satisfying

$$\mathbf{X} \operatorname{vec} U = \operatorname{vec}(XU^T + UX^T), \quad \forall U \in \mathbb{R}^{n \times r}.$$

5.1 Necessary Conditions for the Existence of Spurious Local Minima

As the first step, in the following lemma we obtain necessary conditions for the existence of spurious local minima in the problem (4).

Lemma 13 Assume that the function f in the problem (4) satisfies the δ -RIP_{2r} and κ -BDP_{2r} properties. If X is a local minimizer of (4) and Z is a global minimizer of (4) with $M^* = ZZ^T$, then there exists a symmetric matrix $\mathbf{H} \in \mathbb{R}^{n^2 \times n^2}$ such that the following three conditions hold:

1.
$$\|\mathbf{X}^T \mathbf{H} \mathbf{e}\| \leq 2\kappa \sqrt{\lambda_1(XX^T)} \|\mathbf{e}\|;$$

2. $2I_r \otimes \operatorname{mat}_{\mathrm{S}}(\mathrm{He}) + \mathbf{X}^T \mathbf{H} \mathbf{X} \succeq -2\kappa \|\mathbf{e}\| I_{nr};$

3. **H** satisfies the δ -RIP_{2r} property, i.e., for every matrix $U \in \mathbb{R}^{n \times n}$ with rank $(U) \leq 2r$, it holds that

$$(1-\delta) \|\mathbf{U}\|^2 \le \mathbf{U}^T \mathbf{H} \mathbf{U} \le (1+\delta) \|\mathbf{U}\|^2,$$

where $\mathbf{U} = \operatorname{vec} U$.

Proof Choose **H** to be the matrix satisfying

$$(\operatorname{vec} K)^T \mathbf{H} \operatorname{vec} L = [\nabla^2 f(XX^T)](K, L),$$

for all $K, L \in \mathbb{R}^{n \times n}$. Condition 3 follows immediately from the δ -RIP_{2r} property of the function f. To prove the remaining two conditions, define $g(Y) = f(YY^T)$ and $M = XX^T$. Since X is a local minimizer of the function $g(\cdot)$, for every $U \in \mathbb{R}^{n \times r}$ with $\mathbf{U} = \text{vec } U$, the first-order optimality condition implies that

$$0 = \langle \nabla g(X), U \rangle = \langle \nabla f(M), XU^T + UX^T \rangle.$$
(19)

Define an auxiliary function $h : \mathbb{R}^{n \times n} \to \mathbb{R}$ by letting

$$h(V) = \langle \nabla f(V), XU^T + UX^T \rangle.$$

By the mean value theorem, there exists a matrix ξ on the segment between M and M^* such that

$$[\nabla^2 f(\xi)](M - M^*, XU^T + UX^T) = \langle \nabla h(\xi), M - M^* \rangle = h(M) - h(M^*) = 0, \qquad (20)$$

in which the last equality follows from (19) and $\nabla f(M^*) = 0$. Since rank $(M) \leq r$ and rank $(M^*) \leq r$, we have rank $(\xi) \leq 2r$ and rank $(M - M^*) \leq 2r$. Applying the κ -BDP_{2r} property to the Hessian of $f(\cdot)$ at matrices M and ξ , together with (20), one can obtain

$$\begin{aligned} |\mathbf{e}^{T}\mathbf{H}\mathbf{X}\mathbf{U}| &= |[\nabla^{2}f(M)](M - M^{*}, XU^{T} + UX^{T})| \\ &\leq \kappa \|M - M^{*}\|_{F} \|XU^{T} + UX^{T}\|_{F} \\ &\leq 2\kappa \|\mathbf{e}\| \|XU^{T}\|_{F} \\ &= 2\kappa \|\mathbf{e}\| \sqrt{\operatorname{tr}(UX^{T}XU^{T})} \\ &\leq 2\kappa \|\mathbf{e}\| \sqrt{\lambda_{1}(XX^{T})} \|\mathbf{U}\| \end{aligned}$$

Condition 1 can be proved by setting $\mathbf{U} = \mathbf{X}^T \mathbf{H} \mathbf{e}$.

For every $U \in \mathbb{R}^{n \times r}$ with $\mathbf{U} = \operatorname{vec} U$, the second-order optimality condition gives

$$0 \le [\nabla^2 g(X)](U,U) = [\nabla^2 f(M)](XU^T + UX^T, XU^T + UX^T) + 2\langle \nabla f(M), UU^T \rangle.$$
(21)

The first term on the right-hand side can be equivalently written as $(\mathbf{X}\mathbf{U})^T\mathbf{H}(\mathbf{X}\mathbf{U})$. A similar argument can be made to conclude that there exists another matrix ξ' on the segment

between M and M^* such that

$$\langle \nabla f(M), UU^T \rangle = \langle \nabla f(M) - \nabla f(M^*), UU^T \rangle$$

$$= [\nabla^2 f(\xi')](M - M^*, UU^T)$$

$$\leq [\nabla^2 f(M)](M - M^*, UU^T) + \kappa \|M - M^*\|_F \|UU^T\|_F$$

$$= \operatorname{vec}(UU^T)\mathbf{He} + \kappa \|\mathbf{e}\| \|\mathbf{U}\|^2$$

$$= \frac{1}{2}(\operatorname{vec} U)^T \operatorname{vec}((W + W^T)U) + \kappa \|\mathbf{e}\| \|\mathbf{U}\|^2$$

$$= \mathbf{U}^T (I_r \otimes \operatorname{mat}_S(\mathbf{He}))\mathbf{U} + \kappa \|\mathbf{e}\| \|\mathbf{U}\|^2,$$
(22)

in which $W \in \mathbb{R}^{n \times n}$ is the unique matrix satisfying vec W =**He**. Condition 2 can be obtained by combining (21) and (22).

For given $X, Z \in \mathbb{R}^{n \times r}$ and $\kappa \ge 0$, one can construct an optimization problem based on the conditions in Lemma 13 as follows:

$$\begin{array}{ll} \min_{\delta,\mathbf{H}} & \delta \\ \text{s.t.} & \|\mathbf{X}^T \mathbf{H} \mathbf{e}\| \le a, \\ & 2I_r \otimes \max_{\mathbf{S}}(\mathbf{H} \mathbf{e}) + \mathbf{X}^T \mathbf{H} \mathbf{X} \succeq -bI_{nr}, \\ & \mathbf{H} \text{ is symmetric and satisfies } \delta \text{-RIP}_{2r}, \end{array}$$
(23)

where

$$a = 2\kappa \sqrt{\lambda_1(XX^T)} \|\mathbf{e}\|, \quad b = 2\kappa \|\mathbf{e}\|.$$
(24)

Let $\delta(X, Z; \kappa)$ be the optimal value of (23). Assume that f in the original problem (4) satisfies δ -RIP_{2r} and κ -BDP_{2r}. By Lemma 13, if X is a local minimizer of (4) and Z is a global minimizer of (4) with $M^* = ZZ^T$, then $\delta \geq \delta(X, Z; \kappa)$. As a result, by defining $\delta^*(\kappa)$ as the optimal value of the optimization problem

$$\min_{X,Z \in \mathbb{R}^{n \times r}} \delta(X, Z; \kappa) \quad \text{s.t.} \quad XX^T \neq ZZ^T,$$

the problem (4) is guaranteed to have no spurious local minima as long as $\delta < \delta^*(\kappa)$.

The remaining task is to compute $\delta(X, Z; \kappa)$ and $\delta^*(\kappa)$. First, by the property of the Schur complement, the first constraint in (23) can be equivalently written as

$$\begin{bmatrix} I_{nr} & \mathbf{X}^T \mathbf{H} \mathbf{e} \\ (\mathbf{X}^T \mathbf{H} \mathbf{e})^T & a^2 \end{bmatrix} \succeq 0.$$

The major difficulty of solving (23) comes from the last constraint, since it is NP-hard to verify whether a given quadratic form satisfies δ -RIP_{2r} (Tillmann and Pfetsch, 2014). Instead, we tighten the last constraint of (23) by requiring **H** to have a norm-preserving property for all matrices instead of just for matrices with rank at most 2r, i.e.,

$$(1-\delta) \|\mathbf{U}\|^2 \leq \mathbf{U}^T \mathbf{H} \mathbf{U} \leq (1+\delta) \|\mathbf{U}\|^2, \quad \forall \mathbf{U} \in \mathbb{R}^{n^2},$$

which leads to following semidefinite program:

$$\begin{array}{ll} \min_{\delta,\mathbf{H}} & \delta \\ \text{s. t.} & \begin{bmatrix} I_{nr} & \mathbf{X}^T \mathbf{H} \mathbf{e} \\ (\mathbf{X}^T \mathbf{H} \mathbf{e})^T & a^2 \end{bmatrix} \succeq 0, \\ & 2I_r \otimes \max_{\mathbf{S}}(\mathbf{H} \mathbf{e}) + \mathbf{X}^T \mathbf{H} \mathbf{X} \succeq -bI_{nr}, \\ & (1-\delta)I_{n^2} \preceq \mathbf{H} \preceq (1+\delta)I_{n^2}. \end{array} \tag{25}$$

Similar to the case with linear measurements studied in Zhang et al. (2019), due to the symmetry under orthogonal projections, the problems (23) and (25) turn out to have the same optimal value. This result is a direct generalization of (Zhang et al., 2019, Theorem 8) to the case with nonlinear measurements. See Appendix A for the proof.

Lemma 14 For given $X, Z \in \mathbb{R}^{n \times r}$ and $\kappa \ge 0$, the optimization problems (23) and (25) have the same optimal value.

Even if the value of $\delta(X, Z; \kappa)$ for given X, Z and κ can now be efficiently calculated by solving the semidefinite program (25), to further compute $\delta^*(\kappa)$, an analytical expression is still needed for $\delta(X, Z; \kappa)$. For our purpose, it is sufficient to find a lower bound on $\delta(X, Z; \kappa)$. In the remainder of this section, as a last step to prove Theorem 9 and Theorem 10, we will focus on the problem of lower bounding $\delta(X, Z; \kappa)$ and $\delta^*(\kappa)$. Before digging into this problem, we shall first study the function $\delta(X, Z; \kappa)$ numerically to gain some intuition.

5.2 Numerical Illustration

To numerically analyze $\delta(X, Z; \kappa)$, we select different values for κ , and in each case we sample X and Z randomly by drawing each entry of these matrices independently from the standard normal distribution and then solve the semidefinite program (25) to evaluate $\delta(X, Z; \kappa)$. The empirical cumulative distributions of $\delta(X, Z; \kappa)$ from 5000 samples for n = 5and different κ and rank r are given in Figure 1. It can be observed that when κ increases, $\delta(X, Z; \kappa)$ becomes smaller and a worse bound is expected from finding the minimum value $\delta^*(\kappa)$ of $\delta(X, Z; \kappa)$. On the other hand, $\delta(X, Z; \kappa)$ increases when r grows. For example, in the case when $\kappa = 0.05$, all of the samples satisfy $\delta(X, Z; \kappa) \ge 0.44$ for rank r = 1, $\delta(X, Z; \kappa) \ge 0.51$ for rank r = 2, and $\delta(X, Z; \kappa) \ge 0.64$ for rank r = 3. This observation suggests that the stochastic gradient method may perform better in the higher-rank cases, since its trajectory during the iteration is less likely to be close to a spurious local minimizer, and it will be easier to escape even if the trajectory encounters a spurious solution X that is not detected in the above sampling process.

5.3 Global Guarantee for the Rank-1 Case

When r = 1, X and Z reduce to vectors and henceforth will be denoted as x and z with

$$\mathbf{e} = x \otimes x - z \otimes z, \quad \mathbf{X}u = x \otimes u + u \otimes x, \quad \sqrt{\lambda_1(xx^T)} = \|x\|.$$

Moreover,

$$\|\mathbf{X}u\|^{2} = 2\|x\|^{2}\|u\|^{2} + 2(x^{T}u)^{2}, \quad \forall u \in \mathbb{R}^{n}.$$
(26)

Given two vectors $x, z \in \mathbb{R}^n$ with $x \neq 0$ and $xx^T \neq zz^T$, one can find a unit vector $w \in \mathbb{R}^n$ such that w is orthogonal to x and $z = c_1x + c_2w$ for some scalars c_1 and c_2 . Then,

$$\mathbf{e} = \mathbf{X}\tilde{y} - c_2^2(w \otimes w),$$

in which

$$\tilde{y} = \frac{1 - c_1^2}{2}x - c_1 c_2 w.$$

Note that $\mathbf{X}\tilde{y}$ is orthogonal to $w \otimes w$. Furthermore, since $\tilde{y} \neq 0$ by $xx^T \neq zz^T$ and thus $\mathbf{X}\tilde{y} \neq 0$ by (26), one can rescale \tilde{y} into \hat{y} such that $\|\mathbf{X}\hat{y}\| = 1$ and

$$\mathbf{e} = \|\mathbf{e}\|(\sqrt{1-\alpha^2}\mathbf{X}\hat{y} - \alpha(w \otimes w)), \tag{27}$$

with

$$\alpha := \frac{c_2^2}{\|\mathbf{e}\|} = \frac{\|z\|^2 - (x^T z / \|x\|)^2}{\|\mathbf{e}\|}.$$
(28)

In addition, (26) also implies

$$\|\hat{y}\| \le \frac{\|\mathbf{X}\hat{y}\|}{\sqrt{2}\|x\|} = \frac{1}{\sqrt{2}\|x\|}.$$
(29)

To proceed with our proof, we will need the next lemma that studies the eigenvalues of some structured rank-2 matrices.

Lemma 15 (Zhang et al. (2019)) Let u and v be two vectors of the same dimension. The eigenvalues of the matrix $uv^T + vu^T$ can take only three possible values

$$||u|| ||v|| (1 + \cos \theta), - ||u|| ||v|| (1 - \cos \theta), 0,$$

where θ is the angle between u and v.

Lemma 16 Let $x, z \in \mathbb{R}^n$ with $xx^T \neq zz^T$. The optimal value $\delta(x, z; \kappa)$ of (25) satisfies

$$\delta(x, z; \kappa) \ge \frac{1 - \eta_0(x, z) - 2(1 + \sqrt{2})\kappa}{1 + \eta_0(x, z) + 2(1 + \sqrt{2})\kappa}$$

in which

$$\eta_0(x,z) = \begin{cases} \frac{1-\sqrt{1-\alpha^2}}{1+\sqrt{1-\alpha^2}}, & \text{if } \beta \ge \frac{\alpha}{1+\sqrt{1-\alpha^2}}, \\ \frac{\beta(\beta-\alpha)}{\beta\alpha-1}, & \text{if } \beta \le \frac{\alpha}{1+\sqrt{1-\alpha^2}}, \end{cases}$$

with α defined in $(28)^1$ and $\beta = ||x||^2/||\mathbf{e}||$.

1. When x = 0, α is defined to be $||z||^2/||\mathbf{e}||$.



Figure 1: The empirical probability of $\delta(X, Z; \kappa) \leq \delta_0$ for randomly generated X and Z matrices with n = 5 and different κ and rank r.

Proof Define $\eta(x, z; \kappa)$ to be the optimal value of the following optimization problem:

$$\begin{array}{l} \max_{\eta,\mathbf{H}} & \eta \\ \text{s.t.} & \begin{bmatrix} I_{nr} & \mathbf{X}^T \mathbf{H} \mathbf{e} \\ (\mathbf{X}^T \mathbf{H} \mathbf{e})^T & a^2 \end{bmatrix} \succeq 0, \\ & 2 \operatorname{mat}_{\mathrm{S}}(\mathbf{H} \mathbf{e}) + \mathbf{X}^T \mathbf{X} \succeq -bI_{nr}, \\ & \eta I_{n^2} \preceq \mathbf{H} \preceq I_{n^2}. \end{array} \tag{30}$$

It can be verified that

$$\eta(x, z; \kappa) \ge \frac{1 - \delta(x, z; \kappa)}{1 + \delta(x, z; \kappa)},\tag{31}$$

because given any feasible solution (δ, \mathbf{H}) to (25), the point

$$\left(\frac{1-\delta}{1+\delta}, \frac{1}{1+\delta}\mathbf{H}\right)$$

is also a feasible solution to (30). The reason is that the first and last constraints in (30) naturally hold while the second constraint is satisfied due to

$$2\operatorname{mat}_{S}\left(\frac{1}{1+\delta}\mathbf{H}\mathbf{e}\right) + \mathbf{X}^{T}\mathbf{X} \succeq \frac{1}{1+\delta}(2\operatorname{mat}_{S}(\mathbf{H}\mathbf{e}) + \mathbf{X}^{T}\mathbf{H}\mathbf{X}) \succeq -\frac{b}{1+\delta}I_{nr} \succeq -bI_{nr}.$$

Therefore, to find a lower bound on $\delta(x, z; \kappa)$, we only need to find an upper bound on $\eta(x, z; \kappa)$.

The dual problem of (30) can be written as

$$\min_{U_1, U_2, V, G, \lambda, y} \operatorname{tr}(U_2) + \langle \mathbf{X}^T \mathbf{X} + bI_n, V \rangle + a^2 \lambda + \operatorname{tr}(G),$$
s. t. $\operatorname{tr}(U_1) = 1,$
 $(\mathbf{X}y - v)\mathbf{e}^T + \mathbf{e}(\mathbf{X}y - v)^T = U_1 - U_2,$
 $\begin{bmatrix} G & -y \\ -y^T & \lambda \end{bmatrix} \succeq 0,$
 $U_1 \succeq 0, \quad U_2 \succeq 0, \quad V \succeq 0, \quad v = \operatorname{vec} V.$
(32)

By weak duality, the dual objective value associated with any feasible solution to the dual problem (32) is an upper bound on $\eta(x, z; \kappa)$.

In the case when $x \neq 0$, we fix a constant $\gamma \in [0, \alpha]$ and choose

$$y = \frac{\sqrt{1-\gamma^2}}{\|\mathbf{e}\|}\hat{y}, \quad v = \frac{\gamma}{\|\mathbf{e}\|}(w \otimes w),$$

where \hat{y} and w are the vectors defined before (27). Since $\|\mathbf{X}\hat{y}\| = 1$, $\|w \otimes w\| = 1$ and $\mathbf{X}\hat{y}$ is orthogonal to $w \otimes w$, it holds that

$$\|\mathbf{X}y - v\| = \frac{1}{\|\mathbf{e}\|}.$$

Combined with (27), one can obtain

$$\mathbf{e}^T(\mathbf{X}y - v) = \psi(\gamma),$$

where $\psi(\gamma)$ is given by

$$\psi(\gamma) = \gamma \alpha + \sqrt{1 - \gamma^2} \sqrt{1 - \alpha^2}.$$

Now, define

$$M = (\mathbf{X}y - v)\mathbf{e}^T + \mathbf{e}(\mathbf{X}y - v)^T$$

and decompose

$$M = [M]_{+} - [M]_{-},$$

in which both $[M]_+ \succeq 0$ and $[M]_- \succeq 0$. Let θ be the angle between \mathbf{e} and $\mathbf{X}y - v$. By Lemma 15, it holds that

$$\operatorname{tr}([M]_{+}) = \|\mathbf{e}\| \|\mathbf{X}y - v\| (1 + \cos\theta) = 1 + \psi(\gamma),$$

$$\operatorname{tr}([M]_{-}) = \|\mathbf{e}\| \|\mathbf{X}y - v\| (1 - \cos\theta) = 1 - \psi(\gamma).$$

Then, it is routine to verify that

$$U_1^* = \frac{[M]_+}{\operatorname{tr}([M]_+)}, \quad U_2^* = \frac{[M]_-}{\operatorname{tr}([M]_+)},$$
$$v^* = \frac{v}{\operatorname{tr}([M]_+)}, \quad G^* = \frac{1}{\lambda^*} y^* y^{*T}$$
$$\lambda^* = \frac{\|y^*\|}{a}, \quad y^* = \frac{y}{\operatorname{tr}([M]_+)}$$

forms a feasible solution to the dual problem (32) whose objective value is equal to

$$\frac{\operatorname{tr}([M]_{-}) + \langle \mathbf{X}^T \mathbf{X} + bI_n, V \rangle + 2a \|y\|}{\operatorname{tr}([M]_{+})}.$$
(33)

By (26) and (29), one can write

$$\langle \mathbf{X}^T \mathbf{X} + bI_n, V \rangle = \frac{\gamma}{\|\mathbf{e}\|} (\|\mathbf{X}w\|^2 + b) = \frac{\gamma}{\|\mathbf{e}\|} (2\|x\|^2 + b) = 2(\beta + \kappa)\gamma,$$
(34)

$$2a\|y\| \le \frac{2a\|\hat{y}\|}{\|\mathbf{e}\|} \le 2\sqrt{2}\kappa,\tag{35}$$

where a and b are defined in (24). Substituting (34) and (35) into (33) yields that

$$\eta(x, z; \kappa) \le \Psi(\gamma) + 2(1 + \sqrt{2})\kappa,$$

where

$$\Psi(\gamma) = \frac{2\beta\gamma + 1 - \psi(\gamma)}{1 + \psi(\gamma)}.$$

A simple calculation shows that the function $\Psi(\gamma)$ has at most one stationary point over the interval $(0, \alpha)$ and

$$\min_{0 \le \gamma \le \alpha} \Psi(\gamma) = \eta_0(x, z).$$

In the case when x = 0, we have $\eta_0(x, z) = 0$, and

$$U_1 = \frac{\mathbf{e}\mathbf{e}^T}{\|\mathbf{e}\|^2}, \quad U_2 = 0, \quad V = \frac{zz^T}{2\|\mathbf{e}\|^2}, \\ G = 0, \quad \lambda = 0, \quad y = 0$$

forms a feasible solution to the dual problem (32), which implies that

$$\eta(x, z, \kappa) \le \langle bI_n, V \rangle = \kappa.$$

In either case, it holds that

$$\eta(x, z; \kappa) \le \eta_0(x, z) + 2(1 + \sqrt{2})\kappa,$$

which gives the desired result after combining it with (31).

Proof of Theorem 9 By Lemma 14 and the discussion after Lemma 13, we only need to show that

$$\delta(x, z; \kappa) \ge \frac{2 - 6(1 + \sqrt{2})\kappa}{4 + 6(1 + \sqrt{2})\kappa},\tag{36}$$

for all $x, z \in \mathbb{R}^n$ with $xx^T \neq zz^T$. Similarly to the approach used in proof of (Zhang et al., 2019, Theorem 3), it can be verified that the function $\eta_0(x, z)$ defined in the statement of Lemma 16 has the maximum value 1/3 that is attained by any two vectors x and z that are orthogonal to each other such that ||x||/||z|| = 1/2. Consequently, (36) holds in light of Lemma 16.

5.4 Local Guarantee for the Rank-r Case

The key step in the proof of Theorem 9 is to derive a closed-form expression serving as a lower bound on $\delta(X, Z; \kappa)$. Similar to the idea used in Zhang and Zhang (2020), we need to first simplify $\delta(X, Z; \kappa)$ for the higher-rank cases by removing the constraint corresponding to the second-order optimality condition in (25). The first step is to establish the following lemma that is similar to Lemma 13 but ignores the second-order condition.

Lemma 17 Assume that the function f in the problem (4) satisfies the δ -RIP_{2r} property. If X is a local minimizer of (4) and Z is a global minimizer of (4) with $M^* = ZZ^T$, then there exists a symmetric matrix $\mathbf{H} \in \mathbb{R}^{n^2 \times n^2}$ such that the following two conditions hold:

- 1. $\mathbf{X}^T \mathbf{H} \mathbf{e} = 0;$
- 2. **H** satisfies the δ -RIP_{2r} property, i.e., for every matrix $U \in \mathbb{R}^{n \times n}$ with rank $(U) \leq 2r$, it holds that

$$(1-\delta) \|\mathbf{U}\|^2 \le \mathbf{U}^T \mathbf{H} \mathbf{U} \le (1+\delta) \|\mathbf{U}\|^2,$$

where $\mathbf{U} = \operatorname{vec} U$.

Proof Following the proof of Lemma 13 but using a different **H** that will be given below, we arrive at a matrix $\xi \in \mathbb{R}^{n \times n}$ satisfying rank $(\xi) \leq 2r$ and (20). Choose **H** to be the matrix satisfying

$$(\operatorname{vec} K)^T \mathbf{H} \operatorname{vec} L = [\nabla^2 f(\xi)](K, L),$$

for all $K, L \in \mathbb{R}^{n \times n}$. Then, (20) implies that $\mathbf{e}^T \mathbf{H} \mathbf{X} \mathbf{U} = 0$, which further implies Condition 1 since \mathbf{U} is arbitrary. Condition 2 immediately follows from the δ -RIP_{2r} property of f.

For given $X, Z \in \mathbb{R}^{n \times r}$, one can similarly construct an optimization problem based on the conditions in Lemma 17 as follows:

$$\begin{array}{ll} \min_{\boldsymbol{\delta},\mathbf{H}} & \boldsymbol{\delta} \\ \mathrm{s.\,t.} & \mathbf{X}^T \mathbf{H} \mathbf{e} = \mathbf{0}, \\ & \mathbf{H} \text{ is symmetric and satisfies } \boldsymbol{\delta} \text{-} \mathrm{RIP}_{2r}, \end{array}$$

whose optimal value is the same as that of the following semidefinite program by an argument similar to Lemma 14:

$$\begin{array}{ll}
\min_{\delta,\mathbf{H}} & \delta \\
\text{s. t.} & \mathbf{X}^T \mathbf{H} \mathbf{e} = 0, \\
& (1-\delta) I_{n^2} \preceq \mathbf{H} \preceq (1+\delta) I_{n^2}.
\end{array}$$
(37)

Define $\delta_f(X, Z)$ to be the optimal value of the above problem. Then, the problem (4) has no spurious local minima if the function f satisfies δ -RIP_{2r} such that $\delta < \delta_f(X, Z)$ for all $X, Z \in \mathbb{R}^{n \times r}$ with $XX^T \neq ZZ^T$. Unfortunately, this argument cannot lead to a global guarantee since $\delta_f(0, Z) = 0$ for every $Z \in \mathbb{R}^{n \times r}$ corresponding to $\mathbf{H} = I_{n^2}$. Instead, we will turn to local guarantees on the region of all $X \in \mathbb{R}^{n \times r}$ satisfying

$$\|XX^T - M^*\|_F \le \epsilon \lambda_r(M^*) \tag{38}$$

and prove Theorem 10 by further lower bounding $\delta_f(X, Z)$.

One important difference between the rank-1 and higher-rank cases is that in the latter there are infinitely many matrices $X \in \mathbb{R}^{n \times r}$ that produce the same value for the matrix XX^T . In the proof of Theorem 10, the matrix X is normalized by replacing it with another matrix \tilde{X} with $\tilde{X}^T Z \succeq 0$ while keeping $XX^T = \tilde{X}\tilde{X}^T$. The reason for this normalization operation will be explained in the following lemma, which essentially says that after normalization X and Z cannot be too far away from each other if XX^T and ZZ^T are close.

Lemma 18 (Bhojanapalli et al. (2016b)) Let $X, Z \in \mathbb{R}^{n \times r}$ be two arbitrary matrices such that $Z^T X = X^T Z$ is a positive semidefinite matrix. Then,

$$\lambda_r(ZZ^T) \|Z - X\|_F^2 \le \frac{1}{2(\sqrt{2}-1)} \|ZZ^T - XX^T\|_F^2.$$

Proof of Theorem 10 Let $Z \in \mathbb{R}^{n \times r}$ be a global minimizer of (4) with $ZZ^T = M^*$. To prove by contradiction, assume that there exists a spurious local minimizer $X \in \mathbb{R}^{n \times r}$ satisfying $XX^T \neq ZZ^T$ and (38). Let $X^TZ = UDV^T$ be the singular value decomposition of X^TZ , and define the orthogonal matrix $R = UV^T$. Therefore, the matrix

$$(XR)^T Z = VU^T UDV^T = VDV^T \succeq 0.$$

Furthermore, it is straightforward to verify that XR is also a spurious local minimizer satisfying (38), so $\delta \geq \delta_f(XR, Z)$. On the other hand, applying Lemma 19 (to be stated below) on the local minimizer XR and the global minimizer Z gives rise to the inequality

$$\delta_f(XR, Z) \ge \sqrt{1 - \frac{3 + 2\sqrt{2}}{4}\epsilon^2},$$

which is a contradiction.

Lemma 19 Let $X, Z \in \mathbb{R}^{n \times r}$ such that $M^* = ZZ^T$, $XX^T \neq ZZ^T$, X^TZ is a positive semidefinite matrix, and (38) is satisfied for some $\epsilon \in (0, 2(\sqrt{2} - 1)]$. Then,

$$\delta_f(X,Z) \ge \sqrt{1 - \frac{3 + 2\sqrt{2}}{4}\epsilon^2}.$$

Proof The statement is obviously true when $\lambda_r(M^*) = 0$. If $\lambda_r(M^*) > 0$, by the Wielandt-Hoffman theorem (see Wilkinson (1970)), one can write

$$|\lambda_r(XX^T) - \lambda_r(M^*)| \le ||XX^T - M^*||_F \le \epsilon \lambda_r(M^*),$$

which implies that

$$\lambda_r(XX^T) \ge (1-\epsilon)\lambda_r(M^*) > 0.$$
(39)

Decompose Z as $Z = c_1 X + c_2 W$ for some scalars c_1 and c_2 , where $W \in \mathbb{R}^{n \times r}$ is a matrix satisfying $||W||_F = 1$ and $\langle X, W \rangle = 0$. Then,

$$XX^{T} - ZZ^{T} = (1 - c_{1}^{2})XX^{T} - c_{1}c_{2}(XW^{T} + WX^{T}) - c_{2}^{2}WW^{T}$$

We then choose

$$Y = \frac{1 - c_1^2}{2} X - c_1 c_2 W, \quad y = \operatorname{vec} Y.$$
(40)

Since Y can be written as a linear combination of X and Z while $X^T Z$ is symmetric by assumption, $X^T Y$ is also symmetric and hence $\operatorname{tr}(X^T Y)^2 \geq 0$. Now,

$$\|\mathbf{X}y\|^{2} = \|XY^{T} + YX^{T}\|_{F}^{2}$$

= $2 \operatorname{tr}(X^{T}XY^{T}Y) + \operatorname{tr}(X^{T}Y)^{2} + \operatorname{tr}(Y^{T}X)^{2}$
 $\geq 2 \operatorname{tr}(X^{T}XY^{T}Y)$
 $\geq 2\lambda_{r}(X^{T}X) \operatorname{tr}(Y^{T}Y)$
= $2\lambda_{r}(XX^{T})\|y\|^{2}$.

Note that (40) and $XX^T \neq ZZ^T$ imply that $y \neq 0$, which together with (39) and the above inequality further concludes that $\mathbf{X}y \neq 0$. Moreover,

$$\|\mathbf{e} - \mathbf{X}y\| = \|XX^T - ZZ^T - XY^T - YX^T\|_F = c_2^2 \|WW^T\|_F \le c_2^2,$$

$$\|X - Z\|_F^2 = (1 - c_1)^2 \|X\|_F^2 + c_2^2 \ge c_2^2.$$

Let θ be the angle between **e** and **X**y. It follows from Lemma 18 that

$$\sin \theta \le \frac{\|\mathbf{e} - \mathbf{X}y\|}{\|\mathbf{e}\|} \le \frac{\|X - Z\|_F^2}{\|XX^T - M^*\|_F} \le \frac{\|XX^T - M^*\|_F}{2(\sqrt{2} - 1)\lambda_r(M^*)} \le \frac{\epsilon}{2(\sqrt{2} - 1)} \le 1.$$

Therefore, $\theta < \pi/2$ and

$$\cos\theta \ge \sqrt{1 - \frac{3 + 2\sqrt{2}}{4}\epsilon^2}.$$
(41)

Define $\eta_f(X, Z)$ as the optimal value of the optimization problem

$$\begin{array}{ll} \max_{\eta,\mathbf{H}} & \eta \\ \text{s. t.} & \mathbf{X}^T \mathbf{H} \mathbf{e} = 0, \\ & \eta I_{n^2} \preceq \mathbf{H} \preceq I_{n^2}. \end{array}$$
(42)

Similar to the proof of Lemma 16, it holds that

$$\eta_f(X,Z) \ge \frac{1 - \delta_f(X,Z)}{1 + \delta_f(X,Z)},\tag{43}$$

and it is sufficient to upper bound $\eta_f(X, Z)$ through finding a feasible solution to the dual problem of (42) given by

$$\min_{U_1, U_2, y} \operatorname{tr}(U_2),$$
s. t. $\operatorname{tr}(U_1) = 1,$

$$(\mathbf{X}y)\mathbf{e}^T + \mathbf{e}(\mathbf{X}y)^T = U_1 - U_2,$$

$$U_1 \succeq 0, \quad U_2 \succeq 0.$$
(44)

Let

$$M = (\mathbf{X}y)\mathbf{e}^T + \mathbf{e}(\mathbf{X}y)^T,$$

in which y is defined by (40). The matrix M can be decomposed as

$$M = [M]_{+} - [M]_{-},$$

where $[M]_+ \succeq 0$ and $[M]_- \succeq 0$. By Lemma 15,

$$\operatorname{tr}([M]_{+}) = \|\mathbf{e}\| \|\mathbf{X}y\| (1 + \cos \theta),$$

$$\operatorname{tr}([M]_{-}) = \|\mathbf{e}\| \|\mathbf{X}y\| (1 - \cos \theta).$$

Therefore,

$$U_1^* = \frac{[M]_+}{\operatorname{tr}([M]_+)}, \quad U_2^* = \frac{[M]_-}{\operatorname{tr}([M]_+)}, \quad y^* = \frac{y}{\operatorname{tr}([M]_+)}$$

form a feasible solution to the dual problem (44), which implies that

$$\eta_f(X, Z) \le \frac{\operatorname{tr}([M]_-)}{\operatorname{tr}([M]_+)} = \frac{1 - \cos\theta}{1 + \cos\theta}.$$

This gives the desired result after combining it with (41) and (43).

6. Application: 1-bit Matrix Completion

To demonstrate the effectiveness of the developed conditions for the absence of spurious solutions, in this section we will study the 1-bit matrix completion problem. This is a low-rank matrix recovery problem with nonlinear measurements that naturally arises in applications such as recommendation systems in which each user provides binary (like/dislike) observations (see Davenport et al. (2014); Ghadermarzy et al. (2019)). In this problem, there is an unknown ground truth matrix $M^* \in \mathbb{R}^{n \times n}$ with $M^* \succeq 0$ and rank $(M^*) = r$. One is allowed to take independent measurements on each entry M_{ij}^* , where each measurement value is a binary random variable whose distribution is given by

$$Y_{ij} = \begin{cases} 1 & \text{with probability } \sigma(M_{ij}^*), \\ 0 & \text{with probability } 1 - \sigma(M_{ij}^*). \end{cases}$$

Here, $\sigma(x)$ is commonly chosen to be the sigmoid function $e^x/(e^x + 1)$. Note that M_{ij}^* is an arbitrary real number while the measurements Y_{ij} are restricted to the binary choices 0 and 1. After a large number of measurements are taken, let y_{ij} be the percentage of the measurements on the (i, j)th entry that are equal to 1, for every $i, j \in \{1, \ldots, n\}$. To recover the ground truth matrix, consider its maximum likelihood estimator. The log-likelihood function is given by

$$\sum_{i=1}^{n} \sum_{j=1}^{n} (y_{ij} \log(\sigma(M_{ij})) + (1 - y_{ij}) \log(1 - \sigma(M_{ij}))) = \sum_{i=1}^{n} \sum_{j=1}^{n} (y_{ij} M_{ij} - \log(1 + e^{M_{ij}})).$$

Therefore, the 1-bit matrix completion problem can be formulated as an optimization problem in the form (4) with

$$f(M) = -\sum_{i=1}^{n} \sum_{j=1}^{n} (y_{ij}M_{ij} - \log(1 + e^{M_{ij}})).$$

Our goal is to use Theorem 10 to provide a local guarantee for the absence of spurious local minima for the above problem by showing that there is no spurious local minimizer $X \in \mathbb{R}^{n \times r}$ such that $XX^T \in \overline{\mathcal{B}}(M^*, R)$, where

$$\bar{\mathcal{B}}(M^*, R) = \{ M \in \mathbb{R}^{n \times n} | \| M - M^* \|_F \le R \}.$$

Direct computation shows that

$$[\nabla^2 f(M)](K,L) = \sum_{i=1}^n \sum_{j=1}^n \sigma'(M_{ij}) K_{ij} L_{ij}, \quad \forall M, K, L \in \mathbb{R}^{n \times n}.$$
(45)

For each entry M_{ij}^* , define

$$\underline{M}_{ij}^* = \max\{|M_{ij}^*| - R, 0\}, \quad \overline{M}_{ij}^* = |M_{ij}^*| + R.$$

Since $\sigma'(x)$ is an even function that is decreasing on the region $x \ge 0$, if we let

$$m_1 = \sigma'\left(\min_{i,j} \underline{M}_{ij}^*\right), \quad m_2 = \sigma'\left(\max_{i,j} \overline{M}_{ij}^*\right),$$



Figure 2: The empirical distribution of the radius R of the neighborhood $\overline{\mathcal{B}}(M^*, R)$ that is guaranteed to be free of spurious local minima due to Theorem 10, for randomly generated ground truth matrices M^* with different sizes n and ranks r.

then (45) implies that

$$m_2 \|K\|_F^2 \le [\nabla^2 f(M)](K, K) \le m_1 \|K\|_F^2, \quad \forall M \in \bar{\mathcal{B}}(M^*, R), \ K \in \mathbb{R}^{n \times n}$$

The above inequality shows that $\gamma f(\cdot)$ satisfies δ -RIP_{2r} on the region $\overline{\mathcal{B}}(M^*, R)$ with

$$\gamma = \frac{2}{m_1 + m_2}, \quad \delta = \frac{m_1 - m_2}{m_1 + m_2}$$

Therefore, by Theorem 10, there is no spurious local minimizer X satisfying $XX^T \in \overline{\mathcal{B}}(M^*, R)$ as long as R is sufficiently small to satisfy the inequality

$$\frac{m_1 - m_2}{m_1 + m_2} < \sqrt{1 - \frac{3 + 2\sqrt{2}}{4} \left(\frac{R}{\lambda_r(M^*)}\right)^2}.$$
(46)

In the r = 1 case, if we let

$$m_3 = \max_{i,j} (\sigma'(\underline{M}_{ij}^*) - \sigma'(\overline{M}_{ij}^*))$$

then (45) also implies that

$$\begin{aligned} |[\nabla^2 f(M) - \nabla^2 f(M')](K,L)| &\leq m_3 \sum_{i=1}^n \sum_{j=1}^n |K_{ij}L_{ij}| \\ &\leq m_3 ||K||_F ||L||_F, \quad \forall M, M' \in \bar{\mathcal{B}}(M^*,R), \ K, L \in \mathbb{R}^{n \times n}, \end{aligned}$$

so $\gamma f(\cdot)$ satisfies κ -BDP₂ on the region $\overline{\mathcal{B}}(M^*, R)$ with $\kappa = \gamma m_3$, and then Theorem 9 can be applied similarly to obtain a possibly stronger result.

To illustrate the superiority of our result over previous ones, consider a simple special case in which r = 2 and

$$M^* = \text{diag}(2, 2, 0, \dots, 0)$$

In this case, the inequality (45) becomes

$$\sigma'(2) \|K\|_F^2 \le [\nabla^2 f(M^*)] \le \sigma'(0) \|K\|_F^2,$$

which implies that $\gamma \nabla^2 f(M^*)$ satisfies the δ -RIP₄ property for

$$\gamma = \frac{2}{\sigma'(0) + \sigma'(2)}, \quad \delta = \frac{\sigma'(0) - \sigma'(2)}{\sigma'(0) + \sigma'(2)} \approx 0.41 > 1/5.$$

In addition, it can be observed that the above choice of γ is the best to minimize δ . As a result, the existing bound given in Li et al. (2019) cannot certify the absence of spurious local minima in the region $\bar{\mathcal{B}}(M^*, R)$ no matter how small R is. The reason is that the function $f(\cdot)$ (after scaling) cannot satisfy the δ -RIP₄ property in any local neighborhood of M^* . On the other hand, the above discussion based on our Theorem 10 shows that the problem has no spurious local minima in $\bar{\mathcal{B}}(M^*, R)$ as long as R satisfies (46). Solving the inequality (46) gives R < 1.14.

For an arbitrary ground truth matrix M^* , one can perform a binary search to find the largest R such that the inequality (46) is satisfied and thus conclude that the problem has no spurious local minima in the neighborhood $\overline{\mathcal{B}}(M^*, R)$. For different sizes n and ranks r, Figure 2 plots the empirical distribution of the radius R of such neighborhood for 10^4 random samples of $M^* = ZZ^T$ in which each entry of Z is independently generated from the normal distribution with mean 0 and standard deviation 0.1.

7. Conclusion

In this paper, we first propose the bounded difference property (BDP) in order to study the symmetric low-rank matrix recovery problem with nonlinear measurements. The relationship between BDP and RIP is thoroughly investigated. Then, two novel criteria for the local and global nonexistence of spurious local minima are proposed. It is shown that the developed criteria are superior to the existing conditions relying solely on RIP. In particular, this work offers the first result in the literature on the nonexistence of spurious solutions in a local region for the low-rank matrix recovery problems with nonlinear measurements.

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Appendix A. Proof of Lemma 14

Let OPT(X, Z) denote the optimal value of the optimization problem

$$\begin{array}{ll} \min_{\delta,\mathbf{H}} & \delta \\ \text{s.t.} & \|\mathbf{X}^T \mathbf{H} \mathbf{e}\| \leq a, \\ & 2I_r \otimes \operatorname{mat}_{\mathrm{S}}(\mathbf{H} \mathbf{e}) + \mathbf{X}^T \mathbf{H} \mathbf{X} \succeq -bI_{nr}, \\ & \mathbf{H} \text{ is symmetric and satisfies } \delta \operatorname{-RIP}_{2r}, \end{array} \tag{47}$$

and LMI(X, Z) denote the optimal value of the optimization problem

$$\begin{array}{ll}
\min_{\delta,\mathbf{H}} & \delta \\
\text{s. t.} & \begin{bmatrix} I_{nr} & \mathbf{X}^T \mathbf{H} \mathbf{e} \\ (\mathbf{X}^T \mathbf{H} \mathbf{e})^T & a^2 \end{bmatrix} \succeq 0, \\
& 2I_r \otimes \max_{\mathbf{S}}(\mathbf{H} \mathbf{e}) + \mathbf{X}^T \mathbf{H} \mathbf{X} \succeq -bI_{nr}, \\
& (1-\delta)I_{n^2} \preceq \mathbf{H} \preceq (1+\delta)I_{n^2}.
\end{array}$$
(48)

As mentioned in Section 5, the first constraint in (47) and the first constraint in (48) are interchangeable. Our goal is to prove that OPT(X, Z) = LMI(X, Z) for given $X, Z \in \mathbb{R}^{n \times r}$. Let (v_1, \ldots, v_n) be an orthogonal basis of \mathbb{R}^n such that (v_1, \ldots, v_d) spans the column spaces of both X and Z. Note that $d \leq 2r$. Let $P \in \mathbb{R}^{n \times d}$ be the matrix with the columns (v_1, \ldots, v_d) and $P_{\perp} \in \mathbb{R}^{n \times (n-d)}$ be the matrix with the columns (v_{d+1}, \ldots, v_n) . Then,

$$P^T P = I_d, \quad P_{\perp}^T P_{\perp} = I_{n-d}, \quad P_{\perp}^T P = 0, \quad P^T P_{\perp} = 0$$
$$PP^T + P_{\perp} P_{\perp}^T = I_n, \quad PP^T X = X, \quad PP^T Z = Z.$$

Define $\mathbf{P} = P \otimes P$. Consider the auxiliary optimization problem

$$\begin{array}{ll}
\min_{\delta,\mathbf{H}} & \delta \\
\text{s. t.} & \begin{bmatrix} I_{nr} & \mathbf{X}^{T}\mathbf{H}\mathbf{e} \\ (\mathbf{X}^{T}\mathbf{H}\mathbf{e})^{T} & a^{2} \end{bmatrix} \succeq 0, \\
& 2I_{r} \otimes \max_{\mathbf{S}}(\mathbf{H}\mathbf{e}) + \mathbf{X}^{T}\mathbf{H}\mathbf{X} \succeq -bI_{nr}, \\
& (1-\delta)I_{d^{2}} \preceq \mathbf{P}^{T}\mathbf{H}\mathbf{P} \preceq (1+\delta)I_{d^{2}},
\end{array}$$
(49)

and denote its optimal value as the function $\overline{\text{LMI}}(X, Z)$. Given an arbitrary symmetric matrix $\mathbf{H} \in \mathbb{R}^{n^2 \times n^2}$, if \mathbf{H} satisfies the last constraint in (48), then it obviously satisfies δ -RIP_{2r} and subsequently the last constraint in (47). On the other hand, if H satisfies the last constraint in (47), for every matrix $Y \in \mathbb{R}^{d \times d}$ with $\mathbf{Y} = \text{vec } Y$, since $\text{rank}(PYP^T) \leq d \leq 2r$ and $\text{vec}(PYP^T) = \mathbf{PY}$, by δ -RIP_{2r} property, one arrives at

$$(1-\delta) \|\mathbf{Y}\|^2 = (1-\delta) \|\mathbf{P}\mathbf{Y}\|^2 \le (\mathbf{P}\mathbf{Y})^T \mathbf{H}\mathbf{P}\mathbf{Y} \le (1+\delta) \|\mathbf{P}\mathbf{Y}\|^2 = (1+\delta) \|\mathbf{Y}\|^2,$$

which implies that \mathbf{H} satisfies the last constraint in (49). The above discussion implies that

$$LMI(X, Z) \ge OPT(X, Z) \ge LMI(X, Z).$$

Let

$$\hat{X} = P^T X, \quad \hat{Z} = P^T Z.$$

Lemma 21 and Lemma 22 to be stated later will show that

$$\operatorname{LMI}(X, Z) \le \operatorname{LMI}(\hat{X}, \hat{Z}) \le \overline{\operatorname{LMI}}(X, Z),$$

which completes the proof of Lemma 14.

Before stating Lemma 21 and Lemma 22 that were needed in the proof of Lemma 14, we should first state a preliminary result below.

Lemma 20 Define $\hat{\mathbf{e}}$ and $\hat{\mathbf{X}}$ in the same way as \mathbf{e} and \mathbf{X} , except that X and Z are replaced by \hat{X} and \hat{Z} , respectively. Then, it holds that

$$\mathbf{e} = \mathbf{P}\hat{\mathbf{e}},$$
$$\mathbf{X}(I_r \otimes P) = \mathbf{P}\hat{\mathbf{X}},$$
$$\mathbf{P}^T \mathbf{X} = \hat{\mathbf{X}}(I_r \otimes P)^T.$$

Proof Observe that

$$\mathbf{e} = \operatorname{vec}(XX^T - ZZ^T) = \operatorname{vec}(P(\hat{X}\hat{X}^T - \hat{Z}\hat{Z}^T)P^T) = \mathbf{P}\hat{\mathbf{e}},$$

$$\mathbf{X}(I_r \otimes P)\operatorname{vec}\hat{U} = \mathbf{X}\operatorname{vec}(P\hat{U}) = \operatorname{vec}(X\hat{U}^TP^T + P\hat{U}X^T)$$

$$= \operatorname{vec}(P(\hat{X}\hat{U}^T + \hat{U}\hat{X}^T)P^T) = \mathbf{P}\hat{\mathbf{X}}\operatorname{vec}\hat{U},$$

$$\hat{\mathbf{X}}(I_r \otimes P)^T\operatorname{vec} U = \hat{\mathbf{X}}\operatorname{vec}(P^TU) = \operatorname{vec}(\hat{X}U^TP + P^TU\hat{X}^T)$$

$$= \operatorname{vec}(P^T(XU^T + UX^T)P) = \mathbf{P}^T\mathbf{X}\operatorname{vec} U,$$

where $U \in \mathbb{R}^{n \times r}$ and $\hat{U} \in \mathbb{R}^{d \times r}$ are arbitrary matrices.

Lemma 21 The inequality $\text{LMI}(\hat{X}, \hat{Z}) \ge \text{LMI}(X, Z)$ holds.

Proof Let $(\delta, \hat{\mathbf{H}})$ be an arbitrary feasible solution to the optimization problem defining $\text{LMI}(\hat{X}, \hat{Z})$ with $\delta \leq 1$. It is desirable to show that (δ, \mathbf{H}) with

$$\mathbf{H} = \mathbf{P}\hat{\mathbf{H}}\mathbf{P}^T + (I_{n^2} - \mathbf{P}\mathbf{P}^T)$$

is a feasible solution to the optimization problem defining LMI(X, Z), which directly proves the lemma. To this end, notice that

$$\mathbf{H} - (1-\delta)I_{n^2} = \mathbf{P}(\hat{\mathbf{H}} - (1-\delta)I_{d^2})\mathbf{P}^T + \delta(I_{n^2} - \mathbf{P}\mathbf{P}^T),$$

which is positive semidefinite because

$$\begin{split} I_{n^2} - \mathbf{P}\mathbf{P}^T &= (PP^T + P_{\perp}P_{\perp}^T) \otimes (PP^T + P_{\perp}P_{\perp}^T) - (PP^T) \otimes (PP^T) \\ &= (PP^T) \otimes (P_{\perp}P_{\perp}^T) + (P_{\perp}P_{\perp}^T) \otimes (PP^T) + (P_{\perp}P_{\perp}^T) \otimes (P_{\perp}P_{\perp}^T) \succeq 0. \end{split}$$

Similarly,

$$\mathbf{H} - (1+\delta)I_{n^2} \preceq 0,$$

and therefore the last constraint in (48) is satisfied and \mathbf{H} is always positive semidefinite. Next, since

$$\mathbf{X}^T \mathbf{H} \mathbf{e} = \mathbf{X}^T \mathbf{H} \mathbf{P} \hat{\mathbf{e}} = \mathbf{X}^T \mathbf{P} \hat{\mathbf{H}} \hat{\mathbf{e}} = (I_r \otimes P) \hat{\mathbf{X}}^T \hat{\mathbf{H}} \hat{\mathbf{e}},$$

we have

$$\|\mathbf{X}^T\mathbf{H}\mathbf{e}\|^2 = (\hat{\mathbf{X}}^T\hat{\mathbf{H}}\hat{\mathbf{e}})^T (I_r \otimes P^T) (I_r \otimes P) (\hat{\mathbf{X}}^T\hat{\mathbf{H}}\hat{\mathbf{e}}) = \|\hat{\mathbf{X}}^T\hat{\mathbf{H}}\hat{\mathbf{e}}\|^2,$$

and thus the first constraint in (48) is satisfied. Finally, by letting $W \in \mathbb{R}^{d \times d}$ be the vector satisfying vec $W = \hat{\mathbf{H}}\hat{\mathbf{e}}$, one can write

$$\operatorname{vec}(PWP^T) = \mathbf{P}\operatorname{vec} W = \mathbf{P}\hat{\mathbf{H}}\hat{\mathbf{e}}.$$

Hence,

$$2I_r \otimes \operatorname{mat}_{\mathrm{S}}(\mathbf{H}\mathbf{e}) = 2I_r \otimes \operatorname{mat}_{\mathrm{S}}(\mathbf{H}\mathbf{P}\hat{\mathbf{e}}) = 2I_r \otimes \operatorname{mat}_{\mathrm{S}}(\mathbf{P}\hat{\mathbf{H}}\hat{\mathbf{e}}) = I_r \otimes (P(W + W^T)P^T)$$
$$= 2I_r \otimes (P \operatorname{mat}_{\mathrm{S}}(\hat{\mathbf{H}}\hat{\mathbf{e}})P^T) = 2(I_r \otimes P)(I_r \otimes \operatorname{mat}_{\mathrm{S}}(\hat{\mathbf{H}}\hat{\mathbf{e}}))(I_r \otimes P)^T.$$

In addition,

$$\mathbf{X}^T \mathbf{H} \mathbf{X} (I_r \otimes P) = \mathbf{X}^T \mathbf{H} \mathbf{P} \hat{\mathbf{X}} = \mathbf{X}^T \mathbf{P} \hat{\mathbf{H}} \hat{\mathbf{X}} = (I_r \otimes P) \hat{\mathbf{X}}^T \hat{\mathbf{H}} \hat{\mathbf{X}}.$$

Therefore, by defining

$$\mathbf{S} := 2I_r \otimes \operatorname{mat}_{\mathbf{S}}(\mathbf{He}) + \mathbf{X}^T \mathbf{HX} + bI_{nr},$$

we have

$$(I_r \otimes P)^T \mathbf{S}(I_r \otimes P) = 2I_r \otimes \operatorname{mat}_{\mathbf{S}}(\hat{\mathbf{H}}\hat{\mathbf{e}}) + \hat{\mathbf{X}}^T \hat{\mathbf{H}}\hat{\mathbf{X}} + bI_{dr} \succeq 0,$$

$$(I_r \otimes P_{\perp})^T \mathbf{S}(I_r \otimes P_{\perp}) = (I_r \otimes P_{\perp})^T \mathbf{X}^T \mathbf{H} \mathbf{X}(I_r \otimes P_{\perp}) + bI_{(n-d)r} \succeq 0,$$

$$(I_r \otimes P_{\perp})^T \mathbf{S}(I_r \otimes P) = 0.$$

Since the columns of $I_r \otimes P$ and $I_r \otimes P_{\perp}$ form a basis for \mathbb{R}^{nr} , the above inequalities imply that **S** is positive semidefinite, and thus the second constraint in (48) is satisfied.

Lemma 22 The inequality $\overline{\text{LMI}}(X, Z) \ge \text{LMI}(\hat{X}, \hat{Z})$ holds.

Proof The dual problem of the optimization problem defining $\text{LMI}(\hat{X}, \hat{Z})$ can be expressed as

$$\begin{array}{ll}
\max_{\hat{U}_{1},\hat{U}_{2},\hat{V},\hat{G},\hat{\lambda},\hat{y}} & \operatorname{tr}(\hat{U}_{1}-\hat{U}_{2})-\operatorname{tr}(\hat{G})-a^{2}\hat{\lambda}-b\operatorname{tr}(\hat{V}) \\
\text{s.t.} & \operatorname{tr}(\hat{U}_{1}+\hat{U}_{2})=1, \\
& \sum_{j=1}^{r}(\hat{\mathbf{X}}\hat{y}-\operatorname{vec}\hat{V}_{j,j})\hat{\mathbf{e}}^{T}+\sum_{j=1}^{r}\hat{\mathbf{e}}(\hat{\mathbf{X}}\hat{y}-\operatorname{vec}\hat{V}_{j,j})^{T}-\hat{\mathbf{X}}\hat{V}\hat{\mathbf{X}}^{T}=\hat{U}_{1}-\hat{U}_{2}, \\
& \left[\begin{array}{cc} \hat{G} & -\hat{y} \\ -\hat{y}^{T} & \hat{\lambda} \end{array} \right] \succeq 0, \\
& \hat{U}_{1} \succeq 0, \quad \hat{U}_{2} \succeq 0, \quad \hat{V} = \begin{bmatrix} \hat{V}_{1,1} & \cdots & \hat{V}_{r,1} \\ \vdots & \ddots & \vdots \\ \hat{V}_{r,1}^{T} & \cdots & \hat{V}_{r,r} \end{bmatrix} \succeq 0.
\end{array}$$
(50)

Since

$$\hat{U}_1 = \frac{1}{2d^2}I_{d^2} - \frac{\mu}{2}M, \quad \hat{U}_2 = \frac{1}{2d^2}I_{d^2} + \frac{\mu}{2}M, \quad \hat{V} = \mu I_{dr}, \quad \hat{G} = I_{dr}, \quad \hat{\lambda} = 1, \quad \hat{y} = 0,$$

where

$$M = r((\operatorname{vec} I_d)\hat{\mathbf{e}}^T + \hat{\mathbf{e}}(\operatorname{vec} I_d)^T) + \hat{\mathbf{X}}\hat{\mathbf{X}}^T,$$

is a strict feasible solution to the above dual problem (50) as long as $\mu > 0$ is sufficiently small, Slater's condition implies that strong duality holds for the optimization problem defining $\text{LMI}(\hat{X}, \hat{Z})$. Therefore, we only need to prove that the optimal value of (50) is smaller than or equal to the optimal value of the dual of the optimization problem defining $\overline{\text{LMI}}(X, Z)$ given by:

$$\max_{U_1, U_2, V, G, \lambda, y} \operatorname{tr}(U_1 - U_2) - \operatorname{tr}(G) - a^2 \lambda - b \operatorname{tr}(V)$$
s.t.
$$\operatorname{tr}(U_1 + U_2) = 1,$$

$$\sum_{j=1}^r (\mathbf{X}y - \operatorname{vec} V_{j,j}) \mathbf{e}^T + \sum_{j=1}^r \mathbf{e} (\mathbf{X}y - \operatorname{vec} V_{j,j})^T - \mathbf{X}V\mathbf{X}^T = \mathbf{P}(U_1 - U_2)\mathbf{P}^T,$$

$$\begin{bmatrix} G & -y \\ -y^T & \lambda \end{bmatrix} \succeq 0,$$

$$U_1 \succeq 0, \quad U_2 \succeq 0, \quad V = \begin{bmatrix} V_{1,1} & \cdots & V_{r,1} \\ \vdots & \ddots & \vdots \\ V_{r,1}^T & \cdots & V_{r,r} \end{bmatrix} \succeq 0.$$
(51)

The above claim can be verified by noting that given any feasible solution

$$(\hat{U}_1, \hat{U}_2, \hat{V}, \hat{G}, \hat{\lambda}, \hat{y})$$

to (50), the matrices

$$U_1 = \hat{U}_1, \quad U_2 = \hat{U}_2, \quad V = (I_r \otimes P)\hat{V}(I_r \otimes P)^T,$$
$$\begin{bmatrix} G & -y \\ -y^T & \lambda \end{bmatrix} = \begin{bmatrix} I_r \otimes P & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} \hat{G} & -\hat{y} \\ -\hat{y}^T & \hat{\lambda} \end{bmatrix} \begin{bmatrix} (I_r \otimes P)^T & 0 \\ 0 & 1 \end{bmatrix}$$

form a feasible solution to (51), and both solutions have the same optimal objective value.

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